

# Phase Unwrapping and A Robust Chinese Remainder Theorem

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**Abstract**—In the conventional Chinese remainder theorem (CRT), a small error in a remainder may cause a large error in the solution of an integer, i.e., CRT is not robust. In this letter, we first propose a robust phase unwrapping algorithm with applications in radar signal processing. Motivated from the phase unwrapping algorithm, we then derive a type of robust CRT.

**Index Terms**—Chinese remainder theorem, phase unwrapping, radar signal processing, remainder errors, synthetic aperture radar (SAR).

## I. INTRODUCTION

THE CHINESE remainder theorem (CRT) has tremendous applications in many fields [9], [10], including phase unwrapping [8] in radar signal processing. It is to determine an integer (usually larger) from its remainders (usually smaller) modulo several moduli. The CRT is not robust in the sense that a small error in its remainders may cause a large error in the determined integer by the CRT. CRT with remainder errors has been studied in the literature, see, for example, [1]–[3] and [7] for a generalized CRT, where the remainder errors are corrected by using sufficiently enough remainders/moduli. What is interested in this letter, however, is not necessarily to correct the remainder errors but to reduce the CRT reconstruction errors due to the remainder errors, and there is no requirement on the number of remainders or moduli. To do so, we first propose a robust phase unwrapping algorithm. The problem is described as follows.

Consider a multifrequency antenna array synthetic aperture radar (SAR) [4]. The SAR value at the  $m$ th antenna from the  $i$ th frequency of a moving target is approximately: for  $m = 1, 2, \dots, M$ ,  $i = 1, 2, \dots, L$

$$s_i(m) = c_i \exp\left(j2\pi \frac{\sigma x}{\lambda_i}(m-1)\right) \quad (1)$$

where  $c_i \neq 0$  is independent of  $m$ ,  $M$  is the total number of antennas,  $L$  is the total number of frequencies,  $\lambda_i$  is the wavelength of the transmitted signal with the  $i$ th frequency,  $\sigma$  is a known parameter that is related to the range  $R_0$  of the radar to the moving target and the distance  $d$  between adjacent antennas,  $x$  is an unknown parameter (Doppler shift) to be determined, and  $x$  is related to the speed and location of the moving target.

Manuscript received July 21, 2006; revised August 30, 2006. This work was supported in part by the Air Force Office of Scientific Research (AFOSR) under Grant No. FA9550-05-1-0161 and in part by the National Science Foundation under Grants CCR-0097240 and CCR-0325180. The associate editor coordinating the review of this manuscript and approving it for publication was Prof. Yimin Zhang.

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Digital Object Identifier 10.1109/LSP.2006.884898

Due to the motion of the target, its Doppler shift  $x$  may cause the target shifted in the SAR image domain. Since  $x$  is unknown, the shift amount is unknown, and therefore, a moving target is usually mislocated in a conventional SAR image. However, when  $M > 1$  and  $L > 1$ , a reasonable size  $x$  can be solved, and therefore, a moving target can be correctly relocated in the SAR image [4].

From (1), for each  $i$ ,  $1 \leq i \leq L$ , by taking the  $M$ -point discrete Fourier transform (DFT) of  $s_i(m)$ ,  $1 \leq m \leq M$ , we obtain (only) integer remainders  $k_i$  with  $0 \leq k_i \leq M-1$  as follows:

$$f_i \triangleq \frac{\sigma x}{\lambda_i} = n_i + \frac{k_i}{M} + \epsilon_i \quad (2)$$

where  $n_i$  is a nonnegative unknown integer, and  $\epsilon_i$  is an unknown real number with

$$|\epsilon_i| \leq \frac{1}{2M}. \quad (3)$$

Clearly, to determine the unknown parameter  $x$ , it is sufficient to determine  $f_i$ . Thus, the problem is to determine  $f_i$  from the integer remainders  $k_i$ . From (2), one can see that  $\epsilon_i$  is the precision error, and  $n_i$  is the folding error. While the precision error  $\epsilon_i$  causes a small error on  $f_i$ , the folding error  $n_i$  may cause a large error on  $f_i$ . Thus, in order to robustly determine  $f_i$ , it is necessary to correctly determine  $n_i$ . The problem of interest in this letter is how to uniquely determine  $n_i$ ,  $1 \leq i \leq L$ , in (2) from  $k_i$ ,  $\lambda_i$ ,  $1 \leq i \leq L$ , and  $\sigma$  and  $M$ , where the remainders  $k_i$ ,  $1 \leq i \leq L$  may have errors that may occur when there is an additive noise in signal  $s_i(m)$  in (1). In Section II, we shall propose a solution for this problem under some minor conditions on the wavelengths  $\lambda_i$ .

The above problem is related to the CRT as follows. Let  $\Gamma$  be a positive number such that  $\Gamma_i = \Gamma\lambda_i$ ,  $1 \leq i \leq L$ , are all positive integers. If the precision error  $\epsilon_i = 0$  in (2), then

$$M\Gamma\sigma x = n_i\Gamma_i M + k_i\Gamma_i, \quad 1 \leq i \leq L. \quad (4)$$

Let

$$M_i \triangleq \Gamma_i M, \quad n \triangleq M\Gamma\sigma x, \quad \text{and} \quad r_i \triangleq k_i\Gamma_i, \quad 1 \leq i \leq L \quad (5)$$

then

$$n = n_i M_i + r_i, \quad 0 \leq r_i \leq M_i - 1, \quad 1 \leq i \leq L \quad (6)$$

which is analogous to the CRT problem of determining  $n$  from its remainders  $r_i$  and moduli  $M_i$ ,  $1 \leq i \leq L$ .

## II. ROBUST SOLUTION

When the remainders  $k_i$  in (2) have errors that are possible from the  $M$ -point DFT detection of the signal  $s_i(m)$ , what we have is  $\tilde{k}_i$

$$0 \leq \tilde{k}_i \leq M-1 \quad \text{and} \quad |\tilde{k}_i - k_i| \leq \tau \quad (7)$$

where  $\tau < M$  is the maximal error level in the remainders  $k_i$ . In this case, (2) becomes

$$f_i = \frac{\sigma x}{\lambda_i} = n_i + \frac{\tilde{k}_i}{M} + \tilde{\epsilon}_i, \quad 1 \leq i \leq L \quad (8)$$

where

$$|\tilde{\epsilon}_i| \leq \frac{1+2\tau}{2M}. \quad (9)$$

We want to uniquely determine  $n_i$  in (8) from  $\tilde{k}_i$ ,  $\lambda_i$ , and  $M$ ,  $1 \leq i \leq L$ . To do so, let  $\Gamma$  be the smallest positive number such that

$$\Gamma_i = \Gamma \lambda_i \quad 1 \leq i \leq L \quad (10)$$

are all integers, and  $\Gamma_i$  and  $\Gamma_j$  are co-prime for  $1 \leq i \neq j \leq L$ . Note that this condition may not hold for all possible positive real numbers  $\lambda_i$ ,  $1 \leq i \leq L$ , but there are enough such  $\lambda_i$  in any range. Thus, in what follows, we always assume that the above condition is satisfied. In fact, in SAR applications,  $\lambda_i$  are the wavelengths that we can design by ourselves and therefore flexible to choose.

Without loss of generality, we assume

$$\Gamma_1 < \Gamma_2 < \dots < \Gamma_L \text{ or equivalently } \lambda_1 < \lambda_2 < \dots < \lambda_L. \quad (11)$$

For  $1 \leq i \leq L$ , let

$$\gamma_i \triangleq \Gamma_1 \dots \Gamma_{i-1} \Gamma_{i+1} \dots \Gamma_L \quad (12)$$

where  $\gamma_1 \triangleq \Gamma_2 \dots \Gamma_L$  and  $\gamma_L \triangleq \Gamma_1 \dots \Gamma_{L-1}$ . For each  $i$  with  $2 \leq i \leq L$ , define

$$S_i \triangleq \left\{ (\bar{n}_1, \bar{n}_i) = \underset{\substack{\hat{n}_1 = 0, 1, \dots, \gamma_1 - 1 \\ \hat{n}_i = 0, 1, \dots, \gamma_i - 1}}{\operatorname{argmin}} \left| \hat{n}_i \Gamma_i + \frac{\tilde{k}_i \Gamma_i}{M} - \hat{n}_1 \Gamma_1 - \frac{\tilde{k}_1 \Gamma_1}{M} \right| \right\} \quad (13)$$

and let  $S_{i,1}$  denote the set of all the first components  $\bar{n}_1$  of the pairs  $(\bar{n}_1, \bar{n}_i)$  in set  $S_i$ , i.e.,

$$S_{i,1} \triangleq \{ \bar{n}_1 : (\bar{n}_1, \bar{n}_i) \in S_i \text{ for some } \bar{n}_i \} \quad (14)$$

and define

$$S \triangleq \bigcap_{i=2}^L S_{i,1}. \quad (15)$$

We then have the following result.

*Theorem 1:* Assume  $\Gamma_i$  and  $\Gamma_j$  defined in (10) are co-prime for  $1 \leq i \neq j \leq L$ . If

$$x < \frac{1}{\sigma \Gamma} \Gamma_1 \Gamma_2 \dots \Gamma_L \quad (16)$$

and

$$M > (1 + 2\tau)(\Gamma_1 + \Gamma_L) \quad (17)$$

then set  $S$  defined above contains only element  $n_1$ , i.e.,  $S = \{n_1\}$ , and  $(n_1, \bar{n}_i) \in S_i$  implies  $\bar{n}_i = n_i$  for  $1 \leq i \leq L$ , where  $n_i$ ,  $1 \leq i \leq L$  are the true solution in (8).  $\square$

Its proof is in the Appendix. Note that the case when  $\tau = 0$  has been considered in [4], but there is an error in the result [4, Theorem 1] obtained in [4] where the set  $S_i$  may not necessarily contain only one pair  $(n_1, n_i)$ , and also, the proof in [4] has errors. From the above result, we can see that, when conditions (16) and (17) hold, the folding integers  $n_i$  in (8) can be uniquely solved, and the above results in fact provide an algorithm for the solution of  $n_i$  from the erroneous remainders  $\tilde{k}_i$ . When  $n_i$  in (8) are correctly solved, the unknown parameter  $x$  can be estimated as

$$\hat{x} = \frac{1}{\sigma L} \sum_{i=1}^L \left( n_i + \frac{\tilde{k}_i}{M} \right) \lambda_i \quad (18)$$

and the estimate error can be upper bounded by

$$|x - \hat{x}| \leq \frac{1+2\tau}{2M\sigma} \frac{1}{L} \sum_{i=1}^L \lambda_i. \quad (19)$$

The above estimate error of  $x$  is due to the precision errors  $\epsilon_i$  and the remainder errors  $k_i - \tilde{k}_i$ .

### III. ROBUST CHINESE REMAINDER THEOREM

We now go back to the CRT problem (4)–(6) where the precision errors  $\epsilon_i = 0$  are because the CRT concerns only integers and there are no fractional errors. From (4) and (5), we can see that, for each  $i$  with  $1 \leq i \leq L$ , modulo  $M_i = M\Gamma_i$  and remainder  $r_i$  have a common factor  $\Gamma_i$ , and thus, integer  $n$  also has factor  $\Gamma_i$ . Since a remainder is known *a priori* to have a factor  $\Gamma_i$ , its erroneous version  $\tilde{r}_i$  has a factor  $\Gamma_i$ , too, i.e.,  $\tilde{r}_i = \tilde{k}_i \Gamma_i$ . Thus,

$$\tilde{r}_i - r_i = \epsilon_i \Gamma_i. \quad (20)$$

Assume

$$|\epsilon_i| \leq \tau, \text{ i.e., } |\tilde{k}_i - k_i| \leq \tau. \quad (21)$$

Then, for  $1 \leq i \leq L$

$$n = n_i M_i + \tilde{k}_i \Gamma_i + \epsilon_i \Gamma_i. \quad (22)$$

Since  $n$  has factors  $\Gamma_i$  for  $1 \leq i \leq L$ , integer  $n$  has to have the form  $n = n_0 \Gamma_1 \Gamma_2 \dots \Gamma_L$  for some integer  $n_0$ . Using Theorem 1, we immediately have the following corollary.

*Corollary 1:* Let  $n = n_0 \Gamma_1 \Gamma_2 \dots \Gamma_L$  for some nonnegative integer  $n_0$ . If  $n_0 < M$  and  $M > 2\tau(\Gamma_1 + \Gamma_L)$ , then  $n_i$ ,  $1 \leq i \leq L$ , in (22) can be uniquely determined from  $k_i$ ,  $\Gamma_i$ , and  $M$ ,  $1 \leq i \leq L$ , via (12)–(15). An estimate of  $n$  is

$$\hat{n} = \frac{1}{L} \sum_{i=1}^L (n_i M_i + \tilde{k}_i \Gamma_i) \quad (23)$$

and its error is upper bounded by

$$|\hat{n} - n| = \frac{1}{L} \left| \sum_{i=1}^L \epsilon_i \Gamma_i \right| \leq \frac{\tau}{L} \sum_{i=1}^L \Gamma_i. \quad (24)$$

$\square$

In the conventional CRT, the moduli have to be pair-wisely co-prime, such as  $\Gamma_i$ . In this case, the reconstruction of  $n$  would be expressed as

$$\hat{n} = \sum_{i=1}^L \tilde{r}_i \gamma_i N_i \quad (25)$$

where  $N_i$  is a positive integer such that  $N_i \gamma_i = 1 \pmod{\Gamma_i}$ , and  $\tilde{r}_i$  is an erroneous remainder of  $n$  modulo  $\Gamma_i$  with error  $|\epsilon_i| = |r_i - \tilde{r}_i| \leq \tau$ ,  $1 \leq i \leq L$ . Clearly, the above estimate error due to the erroneous remainders is at least

$$\epsilon_i \Gamma_1 \Gamma_2 \dots \Gamma_{L-1}$$

which is at least in the order of

$$\tau \Gamma_1 \Gamma_2 \dots \Gamma_{L-1}. \quad (26)$$

Comparing (26) with the error estimate (24) in Corollary 1, our proposed robust CRT has much less error than the conventional CRT does. Furthermore, in our case here, the moduli  $M_i = M \Gamma_i$ ,  $1 \leq i \leq L$  and have a common divisor  $M$ , and therefore, the conventional CRT reconstruction formula (25) does not apply. Our study above also provides a reconstruction algorithm of an integer  $n$  from its erroneous remainders when moduli are not pair-wisely co-prime.

#### IV. SIMULATIONS

We now see some simple simulations. We first consider our proposed robust phase unwrapping algorithm in Section II. Consider the case when  $L = 3$ ,  $\lambda_1 = 0.4$ ,  $\lambda_2 = 0.5$ ,  $\lambda_3 = 0.7$ , and  $\sigma = 1$ . We take  $\Gamma = 10$ . Thus,  $\Gamma_1 = 4$ ,  $\Gamma_2 = 5$ , and  $\Gamma_3 = 7$ . In this case, the maximal range of determinable  $x$  in (16) is 14. In our simulations, the unknown parameter  $x$  is uniformly distributed of any real value in the interval  $[0, 14)$ . We take  $M = (1 + 2\tau)(\Gamma_1 + \Gamma_L) + 1$  in (17) and consider the maximal remainder error levels  $\tau = 0, 1, 2, 3, 4, 5$ . The error between the estimate  $\hat{x}$  in (18) and the true value  $x$  is plotted by the solid line marked with  $\times$ , and the error upper bound in (19) is plotted by the solid line marked with  $*$  in Fig. 1. Noticing that the mean value of  $x$  is 7, one can see that the reconstruction errors of  $x$  from the erroneous remainders  $\tilde{k}_i$  in (8) are small. Since the parameter  $x$  is a general real number, there is a precision error  $\epsilon_i$  in (2) that is reflected from error curves shown in Fig. 1 when there is no remainder errors, i.e., when  $\tau = 0$ . The results shown in Fig. 1 are obtained from 200 trials.

For the robust CRT proposed in Section III, we use the same parameters as in Fig. 1, namely,  $L = 3$ ,  $\Gamma_1 = 4$ ,  $\Gamma_2 = 5$ ,  $\Gamma_3 = 7$ , and the maximal remainder error levels  $\tau = 0, 1, 2, 3, 4, 5$ . We also take  $M = (1 + 2\tau)(\Gamma_1 + \Gamma_L) + 1$ . Note that  $M = 2\tau(\Gamma_1 + \Gamma_L) + 1$  is sufficient as stated in Corollary 1 but is not taken to avoid  $M = 1$  when  $\tau = 0$ , which can be, however, easily verified similarly. For an integer  $n = n_0 \Gamma_1 \Gamma_2 \dots \Gamma_L$ , integer  $n_0$  is uniformly taken in the interval  $[0, M - 1]$ , including 0 and  $M - 1$ . The estimate error of  $\hat{n}$  in (23) from the true  $n$  is plotted by the solid line marked with  $\times$  in Fig. 2. The estimate error upper bound in (24) is plotted by the solid line marked with  $*$  in Fig. 2. The errors shown in Fig. 2 can be thought of as absolute errors. The relative errors are shown in Fig. 3, and the curves in Fig. 3 are the error curves shown in Fig. 2 divided by the mean values of  $n$ , respectively. The results shown in Figs. 2 and 3 are obtained from 500 trials. One can see that our robust CRT provides small relative estimate errors when the

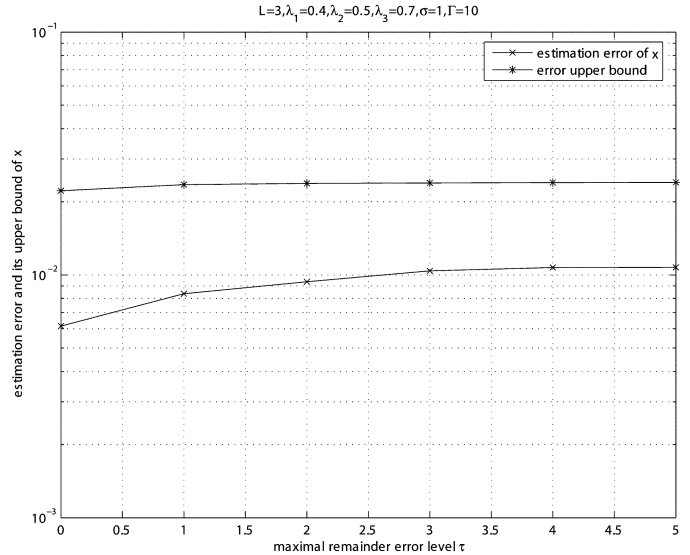


Fig. 1. Estimate error of  $\hat{x}$  in (18) and its upper bound in (19) using proposed robust phase unwrapping algorithm.  $x$  is a uniformly distributed real number in  $[0, 14)$ .

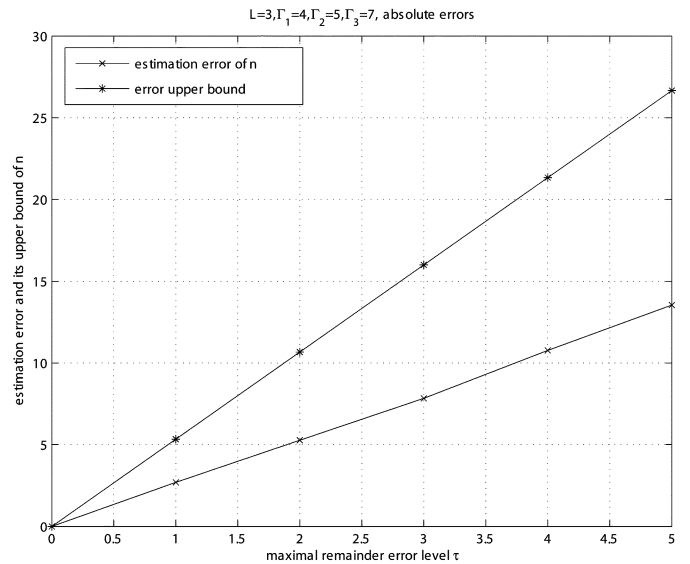


Fig. 2. Absolute estimate error of  $\hat{n}$  in (23) and its upper bound in (24) using proposed robust CRT.

remainders are erroneous, which the conventional CRT cannot provide. Compared to the robust phase unwrapping algorithm dealing with real numbers  $x$  in Fig. 1, we are dealing with integers  $n$  here, and thus, there are no precision errors, i.e.,  $\epsilon_i = 0$ , in (2). The reconstruction (23) is accurate when there is no remainder errors, i.e.,  $\tau = 0$ , which can be verified from Figs. 2 and 3, where the errors are all zero when  $\tau = 0$ .

#### V. CONCLUSION

In this letter, we proposed a robust phase unwrapping algorithm when the remainders have errors. Motivated from the robust phase unwrapping algorithm, we proposed a type of robust CRT for erroneous remainders. As we have mentioned before, what we want to emphasize here is that our proposed robust CRT is not trying to correct remainder errors nor precisely reconstruct an integer as in [1]–[3] but to reduce the reconstruction error

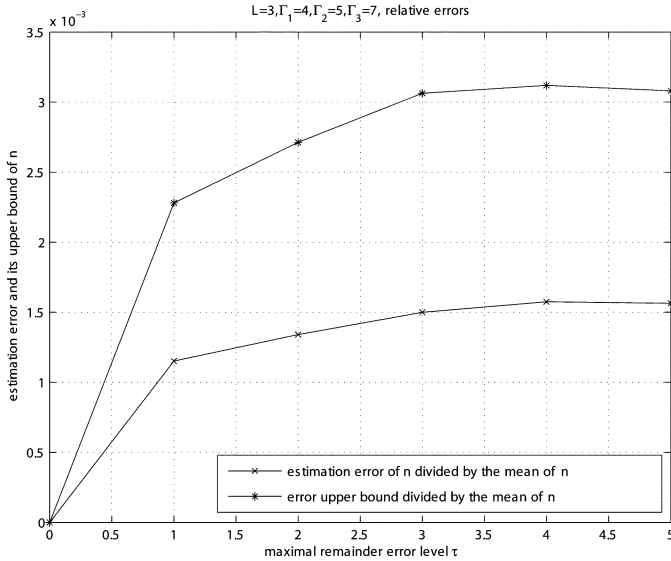


Fig. 3. Relative estimate error of  $\hat{n}$  in (23) and its upper bound in (24) using proposed robust CRT.

from the conventional (non-robust) CRT. Also, what we studied in this letter is only for one integer (one  $x$  or one target) and different from the generalized CRT recently studied in [5]–[7], where multiple integers (multiple  $x$  or multiple targets) and sufficiently many moduli (remainder sets) are involved. Simulation results were provided to verify the theory. As a remark, this letter not only corrected an error on a result obtained in [4] but also considered a more general case for erroneous remainders than in [4].

#### APPENDIX PROOF OF THEOREM 1

From Condition (16) on  $x$ , it is not hard to see that the true solution  $n_i$  in (8) falls in the range  $0 \leq n_i < \gamma_i$  for  $1 \leq i \leq L$ . Thus, for  $2 \leq i \leq L$  and any  $(\bar{n}_1, \bar{n}_i) \in S_i$ , we have

$$\left| \bar{n}_i \Gamma_i + \frac{\tilde{k}_i \Gamma_i}{M} - \bar{n}_1 \Gamma_1 - \frac{\tilde{k}_1 \Gamma_1}{M} \right| \leq \left| n_i \Gamma_i + \frac{\tilde{k}_i \Gamma_i}{M} - n_1 \Gamma_1 - \frac{\tilde{k}_1 \Gamma_1}{M} \right|. \quad (27)$$

From (8)

$$\Gamma \sigma x = n_i \Gamma_i + \frac{\tilde{k}_i \Gamma_i}{M} \Gamma_i + \epsilon_i \Gamma_i, \quad 1 \leq i \leq L. \quad (28)$$

Let  $\mu_i \triangleq \bar{n}_i - n_i$  for  $1 \leq i \leq L$ . From (28), we replace  $(\tilde{k}_i/M)\Gamma_i$  by  $\Gamma \sigma x - n_i \Gamma_i - \epsilon_i \Gamma_i$  in both sides of (27) and have

$$|\mu_i \Gamma_i - \mu_1 \Gamma_1 - (\epsilon_i \Gamma_i - \epsilon_1 \Gamma_1)| \leq |\epsilon_i \Gamma_i - \epsilon_1 \Gamma_1|.$$

Thus, using (9) and (17), we have

$$|\mu_i \Gamma_i - \mu_1 \Gamma_1| \leq 2|\epsilon_i \Gamma_i - \epsilon_1 \Gamma_1| \leq \frac{1+2\tau}{M}(\Gamma_1 + \Gamma_i) < 1. \quad (29)$$

Since  $\mu_i$ ,  $\Gamma_i$ ,  $\mu_1$ , and  $\Gamma_1$  are all integers, (29) implies

$$\mu_i \Gamma_i = \mu_1 \Gamma_1, \quad i = 2, 3, \dots, L. \quad (30)$$

Since  $\Gamma_i$  and  $\Gamma_1$  are co-prime, (30) implies

$$\begin{aligned} \mu_1 &= m_i \Gamma_i \text{ and } \mu_i = m_i \Gamma_1 \\ \text{i.e., } \bar{n}_1 &= n_1 + m_i \Gamma_i \text{ and } \bar{n}_i = n_i + m_i \Gamma_1 \end{aligned} \quad (31)$$

for some integers  $m_i$  with  $|m_i| < \min(\gamma_i, \gamma_1)$ . Replacing (31) into (27), we find that

$$\begin{aligned} \left| n_i \Gamma_i + \frac{\tilde{k}_i \Gamma_i}{M} - n_1 \Gamma_1 - \frac{\tilde{k}_1 \Gamma_1}{M} \right| \\ = \left| \bar{n}_i \Gamma_i + \frac{\tilde{k}_i \Gamma_i}{M} - \bar{n}_1 \Gamma_1 - \frac{\tilde{k}_1 \Gamma_1}{M} \right| \end{aligned} \quad (32)$$

which means  $(n_1, n_i) \in S_i$  for  $i = 2, 3, \dots, L$ . This proves  $n_1 \in S$ . We next show  $S = \{n_1\}$ . Property (31) also implies

$$S_i = \{(n_1 + m_i \Gamma_i, n_i + m_i \Gamma_1) : \text{for some integers } m_i \text{ with } |m_i| < \min(\gamma_i, \gamma_1)\}. \quad (33)$$

If  $\bar{n}_1 \in S$ , then  $\bar{n}_1 \in S_{i,1}$  for  $i = 2, 3, \dots, L$ , and therefore, from the definition of  $S_{i,1}$  in (14) and (33), we have  $\bar{n}_1 - n_1 = m_i \Gamma_i$  for some integer  $m_i$  with  $|m_i| < \min(\gamma_i, \gamma_1)$  for  $i = 2, 3, \dots, L$ . This implies that  $\bar{n}_1 - n_1$  divides all  $\Gamma_i$  for  $i = 2, 3, \dots, L$ , and therefore, from (12),  $\bar{n}_1 - n_1$  is a multiple of  $\gamma_1$ . Since  $0 \leq \bar{n}_1, n_1 \leq \gamma_1 - 1$ , we conclude  $\bar{n}_1 - n_1 = 0$ . This proves that  $S = \{n_1\}$ . In the meantime,  $\bar{n}_1 = n_1$  implies  $m_i = 0$  in (33), i.e.,  $\bar{n}_i = n_i$  for  $i = 2, 3, \dots, L$ . Hence, Theorem 1 is proved. As a remark, despite the fact that [4] only considers the case when  $\tau = 0$ , the proof in [4] has errors.

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