

A literal is either an atomic formula (positive literal) or the negation of an atomic formula (negative literal). A clause contains literals and is interpreted as their disjunction. If a clause contains the literals  $l_1, \dots, l_k$  then we will write it as  $[l_1, \dots, l_k]$ , which is *different* from the notation used in the textbook. Thus,  $[l_1, \dots, l_k]$  is interpreted the same way as the formula  $(l_1 \vee \dots \vee l_k)$ .

Here is a (*non-deterministic*) algorithm to convert a formula,  $A$ , into a set of clauses. Start with  $\{[A]\}$ .

At any point, we will have a set that has the form,  $\{C_1, \dots, C_i, \dots, C_k\}$ , where the  $C$ 's will be in the form  $[A_1, \dots, A_n]$ , and the  $A$ 's are formulae. We are finished with a  $C$  if it contains only literals. (i.e.,  $C$  is already a clause)

While some  $C$  has a non-literal

Let  $C_i$  include a non-literal.

WLOG, we can express  $C_i$  as  $[A_1, \dots, A_{n-1}, A_n]$  where  $A_n$  is a non-literal.

Case 1: (double negation case)

$A_n = \neg\neg B$ , for some formula  $B$ .

Replace  $C_i$  by  $[A_1, \dots, A_{n-1}, B]$ .

Case 2: (disjunctive case).

$A_n$  is a "disjunctive" formula, say  $\beta$ .

Replace  $C_i$  by  $[A_1, \dots, A_{n-1}, \beta_1, \beta_2]$ , where

$\beta_1$  and  $\beta_2$  can be determined from  $\beta$  by using the table in Page 18 of the textbook.

Case 3: (conjunctive case).

$A_n$  is a "conjunctive" formula, say  $\alpha$ .

Replace  $C_i$  by  $C_i^1$  and  $C_i^2$  where

$C_i^1 = [A_1, \dots, A_{n-1}, \alpha_1]$ ,

$C_i^2 = [A_1, \dots, A_{n-1}, \alpha_2]$ , and

$\alpha_1$  and  $\alpha_2$  can be determined from  $\alpha$  by using the table in Page 18 of the textbook.

**Example:** Converting  $A = (\neg p \vee q) \rightarrow (r \vee p)$  into a set of clauses:

$\{[(\neg p \vee q) \rightarrow (r \vee p)]\}$  (starting with  $\{[A]\}$ )

$\{[\neg(\neg p \vee q), (r \vee p)]\}$  (applying Rule 2)

$\{[\neg(\neg p \vee q), r, p]\}$  (applying Rule 2)

$\{[\neg\neg p, r, p], [\neg q, r, p]\}$  (applying Rule 3)

$\{[p, r], [p \neg q, r]\}$  (applying Rule 1)

Note we write  $[p, r]$  rather than  $[p, r, p]$  and that  $[\neg q, r, p]$  can also be written as  $[p, \neg q, r]$ .