

Overview Roadmap

- **Language Translators: Interpreters & Compilers**
- **Context of a compiler**
- **Phases of a compiler**
- **Compiler Construction tools**
- **Terminology**
- **How related to other CS**
- **Goals of a good compiler**

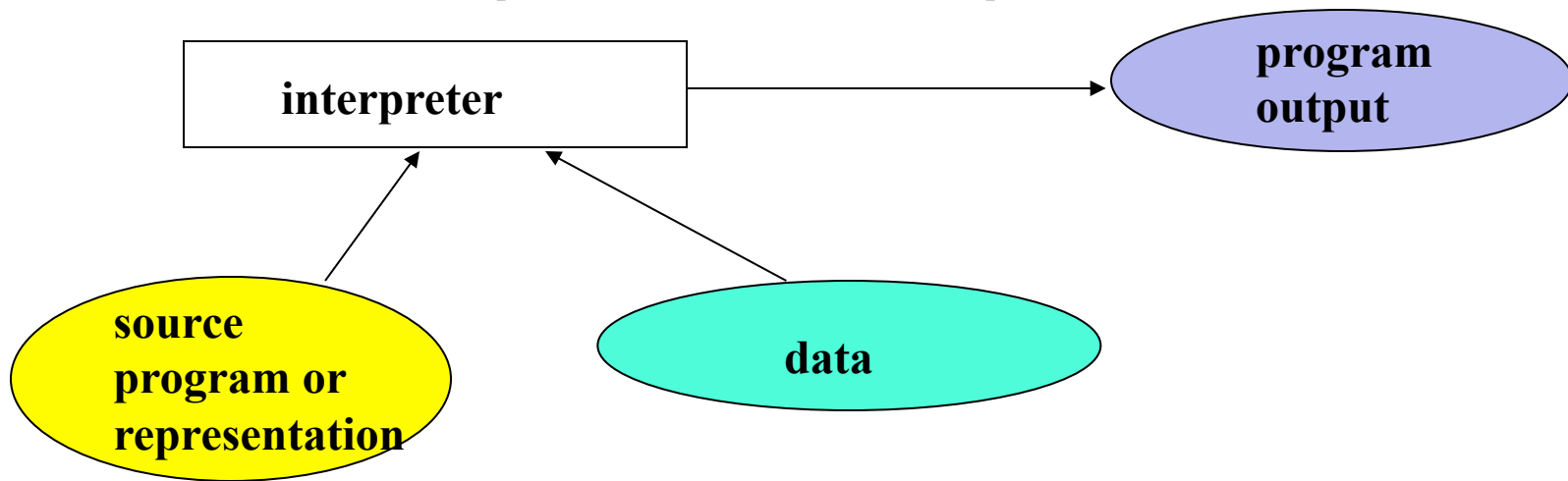
Compilers and Interpreters

- **What is a compiler?**

- **What is an interpreter?**

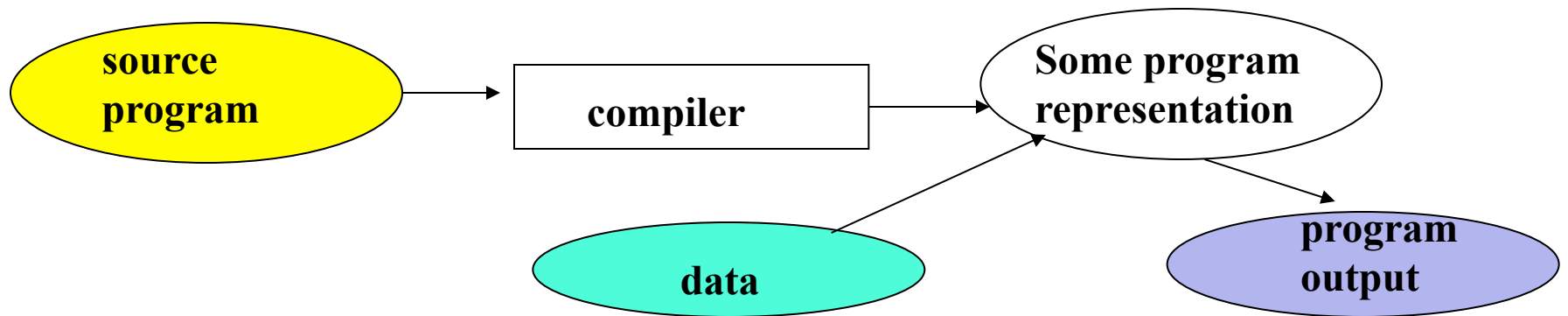
- **Implementation of Languages**

Overview of interpreters and compilers



Locus of control - in interpreter, not program

Compiler has distinct translation and execution phase



Interpreters

Advantages

Disadvantages

Compilers

Advantages

Disadvantages

Relation to Other CS

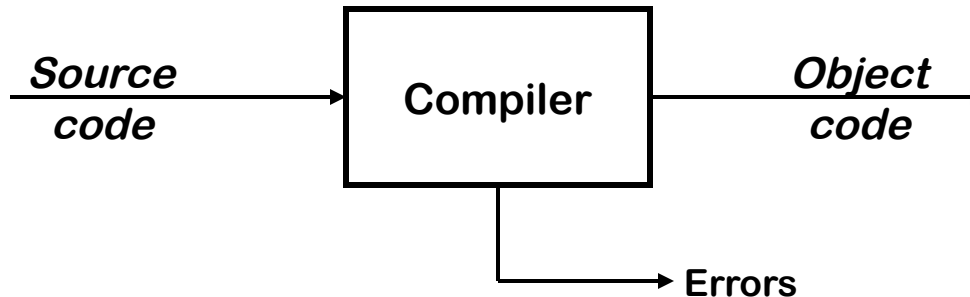
| | |
|---------------------------------------|--|
| <i>Artificial intelligence</i> | Greedy algorithms, Genetic algorithms Heuristic search techniques |
| <i>Algorithms</i> | Graph algorithms, union-find Dynamic programming |
| <i>Theory</i> | DFAs & PDAs, pattern matching Fixed-point algorithms |
| <i>Systems</i> | Allocation & naming, Synchronization, locality |
| <i>Architecture</i> | Memory hierarchy management Functional units & pipelines Instruction set use |

From Your Experience

- You have used several compilers.
- What qualities do you want in a compiler that you buy ?



High-level View of a Compiler



Must recognize legal (and illegal) programs

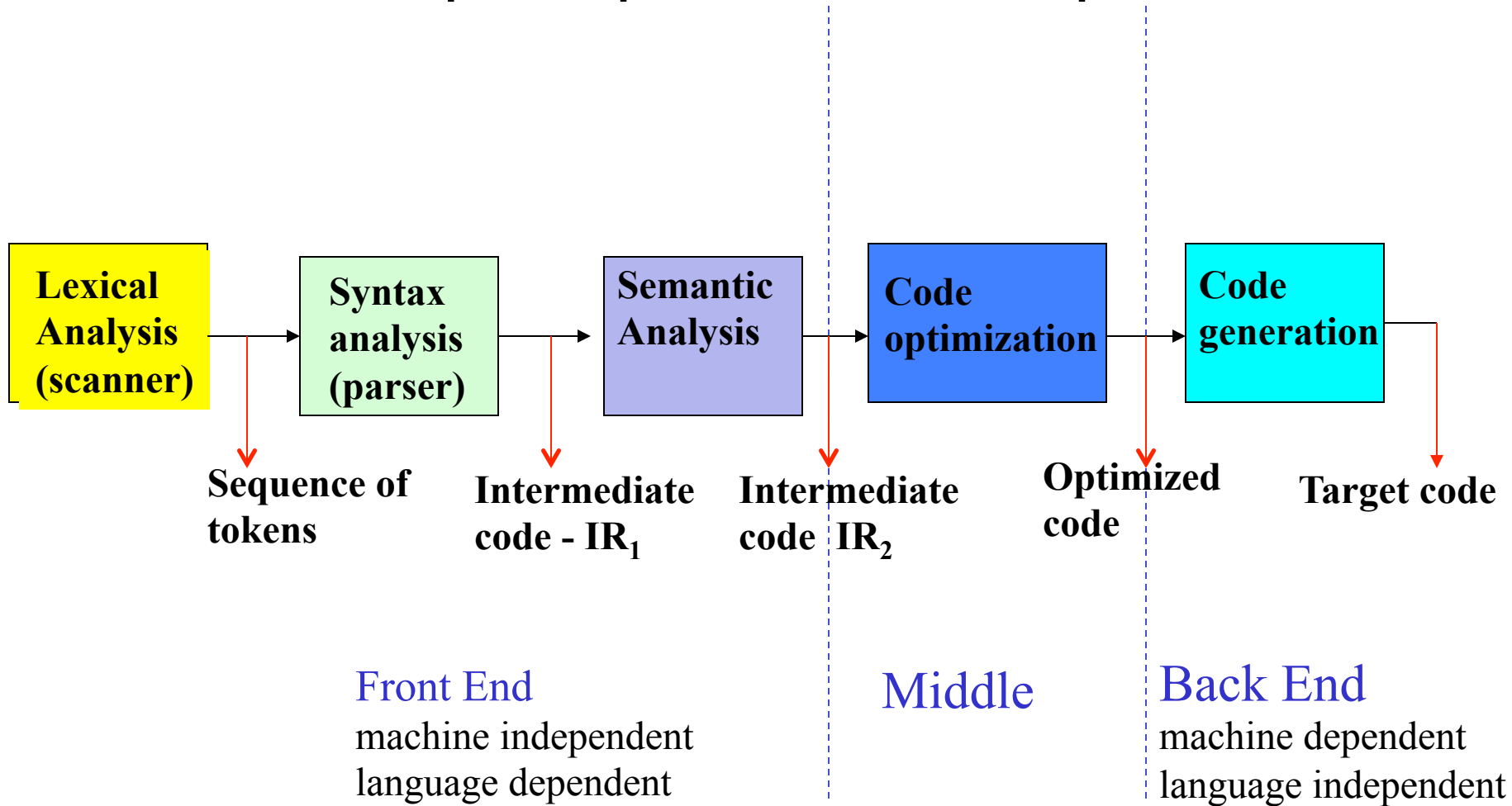
Must generate correct code

Must manage storage of all variables (and code)

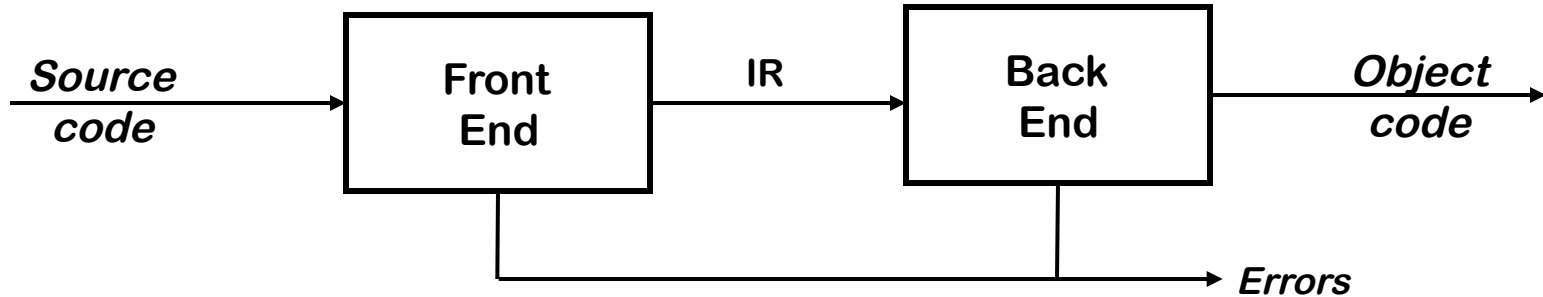
Must agree with OS & linker on format for object code

Big step up from assembly language—use higher level notations

Conceptual phases of compiler



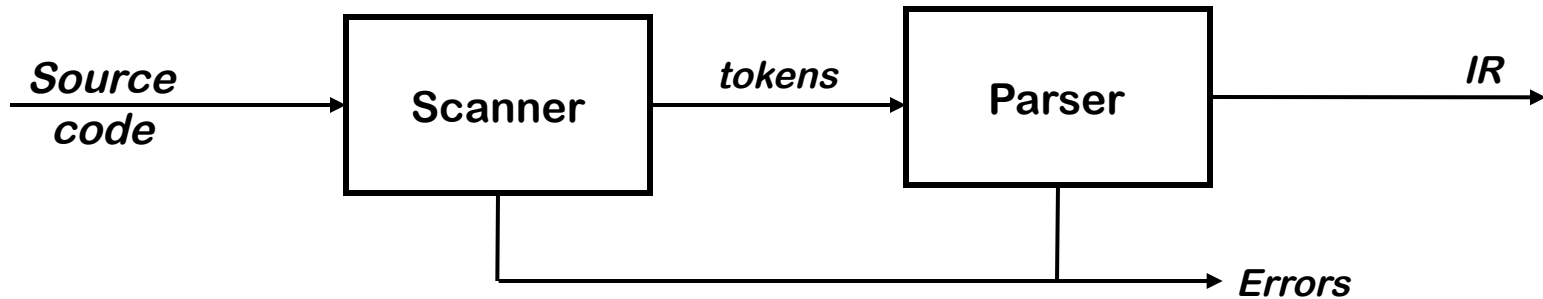
Traditional Two-pass Compiler



Allow 2 passes:

- Use an intermediate representation (IR)
- Front end maps legal source code into IR
- Back end maps IR into target machine code
- Admits multiple front ends & multiple passes (*better code*)

The Front End



Responsibilities

- Recognize legal (& illegal) programs
- Report errors in a useful way
- Produce IR & preliminary storage map
- Shape the code for the back end
- Much of front end construction can be automated

Lexical Analysis/Scanner

Purpose: recognize words - smallest unit

Analyze string of characters from source - left to right to recognize units

Character string - lexeme

Type of lexical entity - token

Smallest unit above letters

Example:

Max:= initial * late + 60

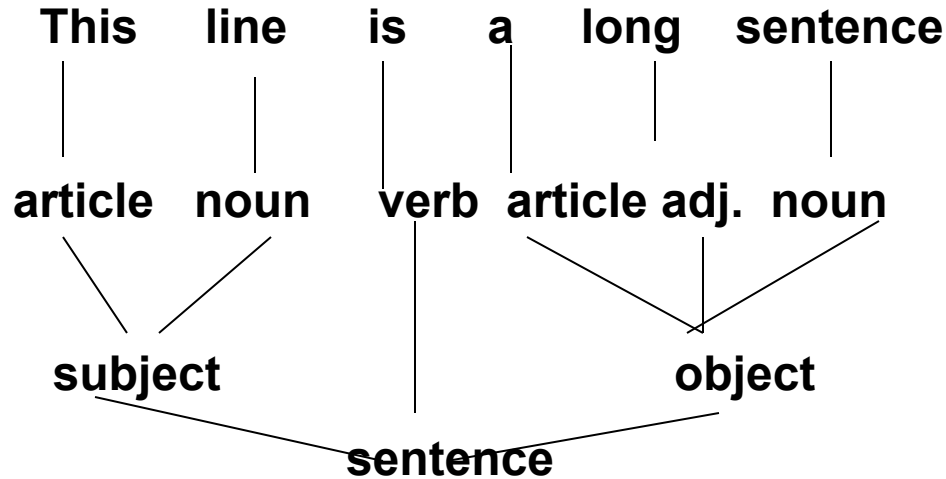
Lexemes: “max”, “:=”, “initial”, “*”, “late”, “+”, “60”

Tokens: Id Id Id := * + Int

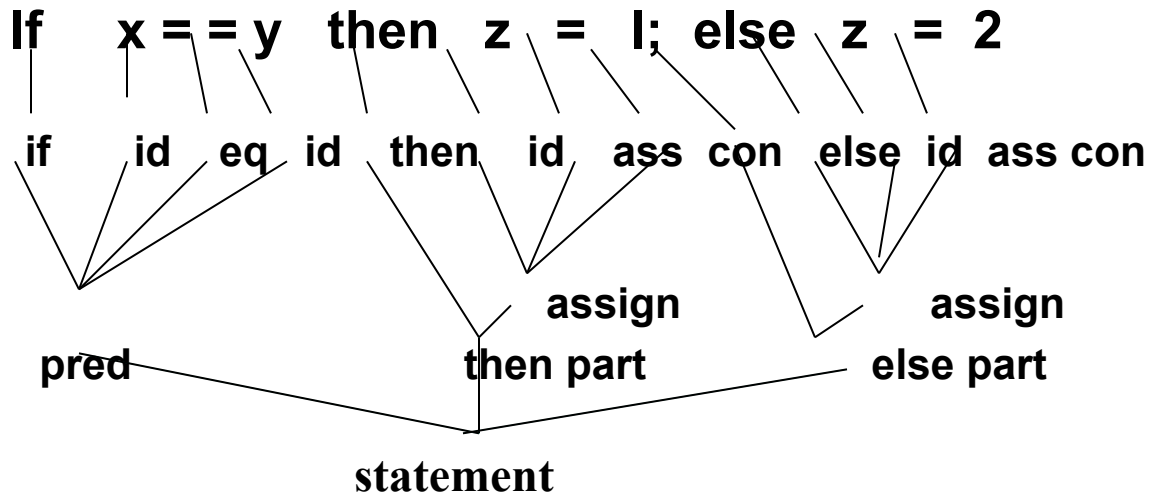
Must recognize blanks, other characters such as % , \$, etc

Syntax Analyzer - Parser

Parsing similar to diagramming a natural language sentence



Parsing



Semantic Analysis

Once structure is understood, determine the meaning using the structure.

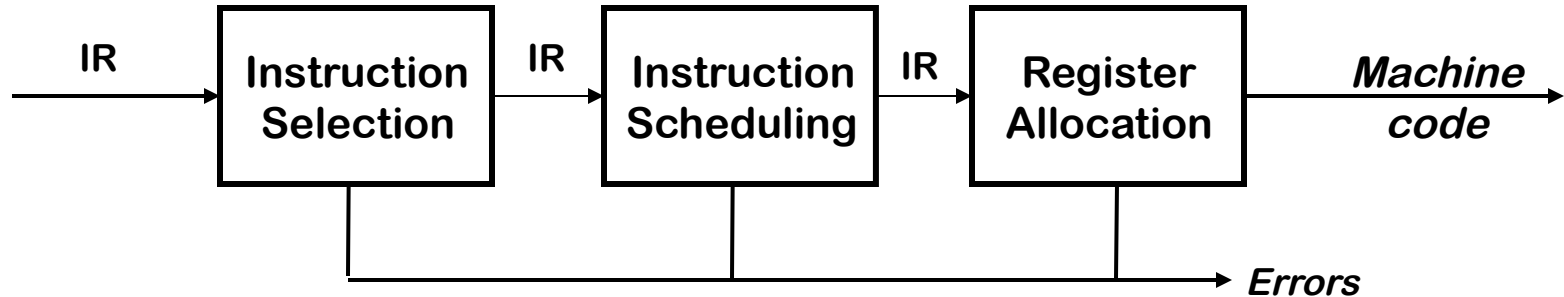
Checks performed to ensure components fit together meaningfully

- information is added to structures
- limited analysis to catch inconsistencies - e.g., type checking

Put semantic meaning in structure -

- produce intermediate form - IR - many forms of IR
- easier to generate machine code from IR
- can be different levels of IR - descending levels of abstraction
 - Highest is source
 - Lowest is target code

The Back End/Code Generation

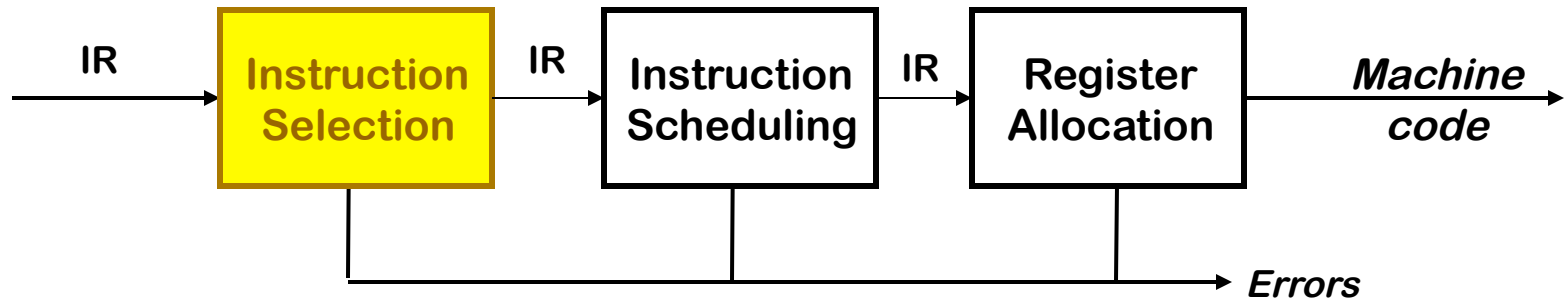


Responsibilities

- Translate IR into target machine code
- Choose instructions to implement each IR operation
- Decide which value to keep in registers
- Ensure conformance with system interfaces

Automation has been much less successful in the back end

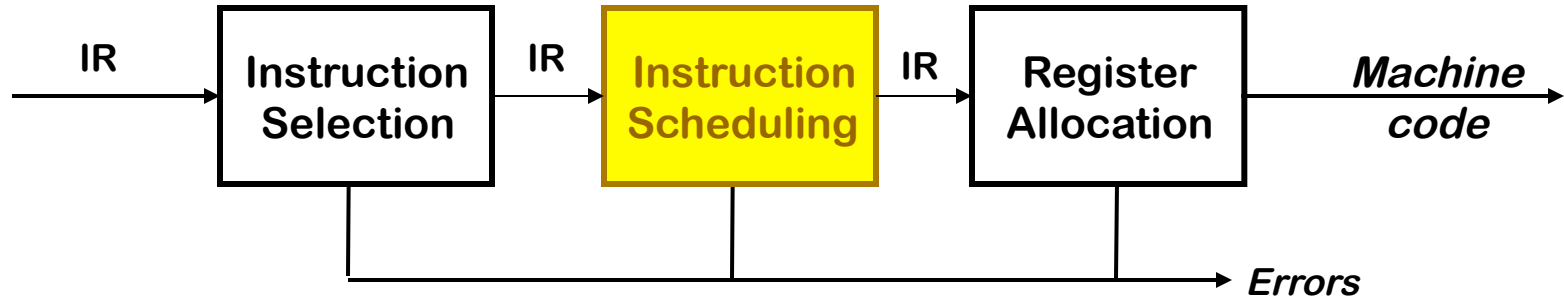
The Back End



Instruction Selection

- Produce fast, compact code
- Take advantage of target features such as addressing modes
- Usually viewed as a pattern matching problem
 - ad hoc methods, pattern matching, dynamic programming
 - Depends on architecture - CISC, RISC

The Back End

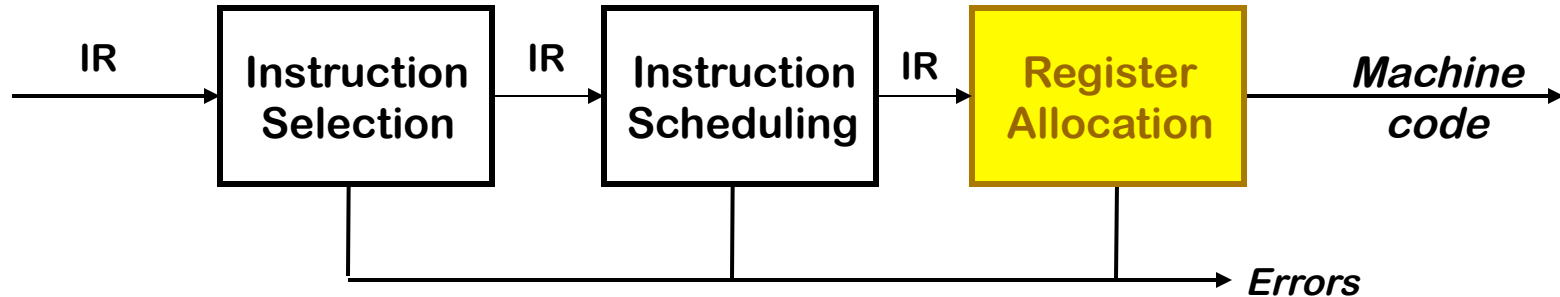


Instruction Scheduling

- Avoid hardware stalls and interlocks
- Use all functional units productively
- Can increase lifetime of variables (**changing the allocation**)
- Optimal scheduling is NP-Complete in nearly all cases

Good heuristic techniques are well understood

The Back End



Register allocation

- Have each value in a register when it is used
- Manage a limited set of resources
- Can change instruction choices & insert LOADs & STOREs
- Optimal allocation is NP-Complete
- Compilers approximate solutions to NP-Complete problems

Code Generation – what kind of code

Produce target code - various forms of target code

1. Assembly Code - symbolic instruction and addresses

- **Easier but not done in modern compilers - assembler slow**

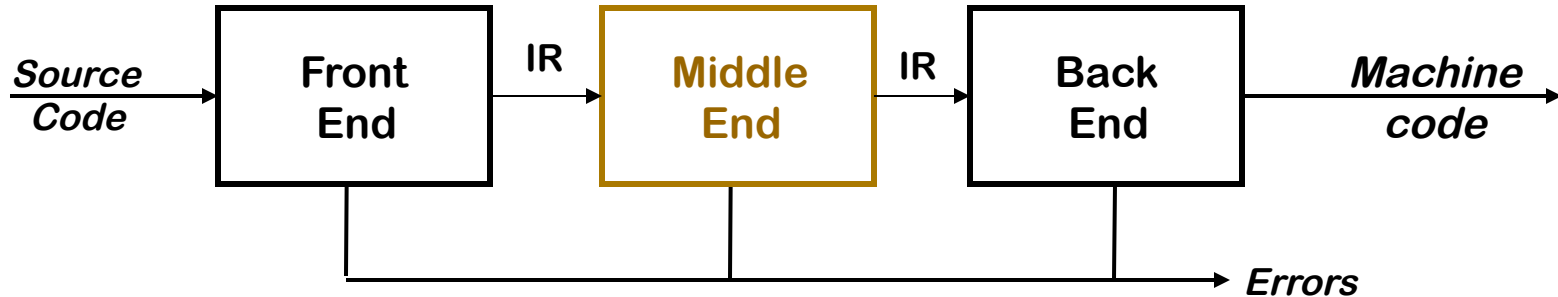
2. Relocatable format –

- **Binary form except external references, instruction addresses and data addresses not bound to address**
- **Need linker and loader**

Both assembly & relocatable allow program modules to be separately compiled

3. Another language

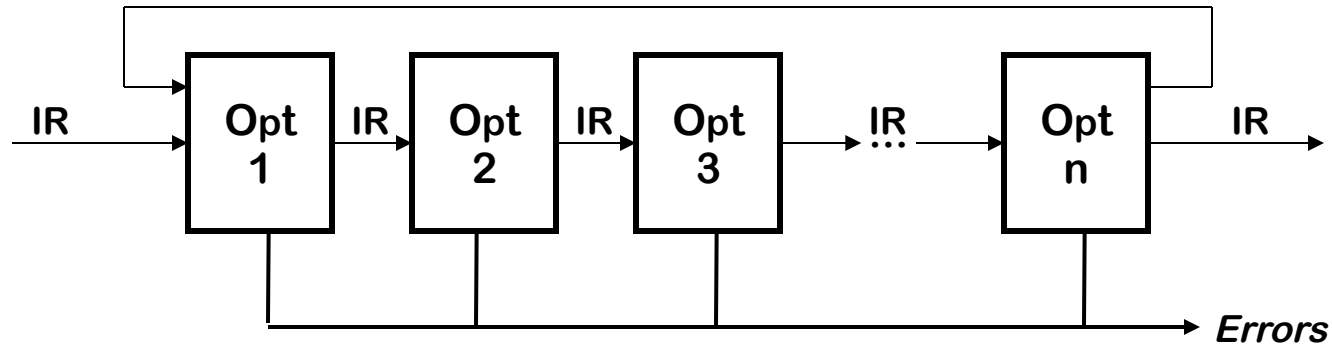
Traditional Three-pass Compiler



Code Improvement (or Optimization)

- Analyzes IR and rewrites (or transforms) IR
- Primary goal is to reduce running time of the compiled code
 - May also improve space, power consumption, ...
- Must preserve “meaning” of the code
 - Measured by values of named variables

The Optimizer (or Middle End)



Modern optimizers are structured as a series of passes

Typical Transformations

- Discover & propagate some constant value
- Move a computation to a less frequently executed place
- Specialize some computation based on context
- Discover a redundant computation & remove it
- Remove useless or unreachable code
- Encode an idiom in some particularly efficient form

Code Optimization

Modify program representation so that the program

- runs faster
- uses less memory
- uses less power
- in general, reduce the resources consumed

e.g., constant propagation and folding

$Y := 3$

$X := Y + 4$

optimizes to $X := 7$

Symbol Table Manager

Collect and maintain information about id's

- attributes e.g., storage allocation, type, scope, number and type of parameters

Usually cuts across all phases - lexical, parsing and semantic, code optimization, code generation

- Phase add information - lexical, parsing and semantic
- Phases use information - code optimization, code generation

Debuggers uses some form of symbol table

Error Reporting

Phases deal with errors - 1st 3 phases handled bulk of errors

Lots of success here

Distinction between phases and passes

Passes - number of times through a program representation

- 1 - passes, 2 - passes, multiple passes
- Languages become more complex - more passes

Phases - conceptual and sometimes physical stages

- Symbol table coordinating information between phases

However, phases are not completely separate - semantic phase must do things that syntax phase should do if it could

Some interaction possible:

- optimization and code generation - what optimizer does affects code generator

Compiler tools

Scanner generator

- Generate lexical analyzer from specification of tokens based on regular expressions
- Examples: Lex, Flex, JLex

Parser generator

- Generate parser from specification of syntactical structure using BNF grammars
- Example: YACC, Bison, CUP

What about compiler generator?

- How do you specify semantics that is useful for compiler?
- How do you specify the architecture?
- How do you specify optimizations?