CISC 320 010 Introduction to Algorithms (Fall 2005)

Homework 1 Handed out: September 8, 2005 **Due date: September 20, 2005**

Your handwriting must be legible, and your answers should be rigorous, concise and in proper order. Please note that the work handed in must be your own.

1. CLRS 3-1 (18 points)

2. (30 points) Indicate, for each pair of expressions (A, B) in the table below, whether A is O, o, Ω , or Θ of B. Assume that are c > 1 and ε >0 constants. Your answer should be in the form of the table with "yes" or "no" written in each box.

	А	В	0	0	Ω	Θ
1	√n	n ^{sin n}				
2	n ^{lg(c)}	c ^{lg(n)}				
3	2^n	$n^{n/2}$				
4	lg(n)	n ^ε				
5	lg(n!)	$lg(n^n)$				

3. CLRS 4.3-2 (17 points) The recurrence $T(n) = 7T(n/2) + n^2$ describes the running time of an algorithm *A*. A competing algorithm *A*' has a running time of $T'(n) = aT'(n/2) + n^2$. What is the largest integer value for *a* such that *A*' is asymptotically faster than *A*?

4. (35 points) Algorithm A is a divide and conquer type with recurrence relation $T(n) = 4T(n/2) + n^2 lg(n)$. Answer the following questions about the recursive tree of A on input size $n = 2^m$, where m is an positive integer.

- a. The root starts at level 0, then what is the input size to a node at level *i*?
- b. How many nodes at level *i*?
- c. How much total work is actually done at level *i*?
- d. The root is at level 0. What level are the leaves at?
- e. Write a series that represent the running time for the algorithm for inputs of size $n = 2^m$, and give explicitly the starting and ending value for the index in the series.
- f. Solve the series.
- g. Give the tightest *big-O* bound for T(n).