- path lengths
 - Intuition: avg number of nodes s steps away from random i is c^s . We reach all vertices when $c^s \approx n$, i.e., $s \approx \ln n / \ln c$.
 - Problem: this argument doesn't work when s is large.
 - Consider two starting vertices i and j with their s- and t-distance neighborhoods, respectively, when s, t are small
 - 1. if "- - " exists between surfaces then one can show that there are edges between larger surfaces



 \implies $\Pr[d_{ij} > s + t + 1] \approx \operatorname{prob} \not\exists$ edge between two surfaces c^s , and c^t when t is small

2. There are on avg $c^s \times c^t$ pairs of nodes, s.t. one lies on each surface and each pair is connected with prob p = c/(n-1)i.e., $\Pr[d_{ij} > s + t + 1] = (1-p)^{c^{s+t}} = (1-c/n)^{c^{l-1}}$ or $\ln \Pr[d_{ij} > l] = c^{l-1} \ln(1-c/n) \approx -c^l/n$

l = *s*+*t*+1

Networks	# of nodes	Diameter					Modularity				
		Observed	Expected	% difference	Z- score ¹⁴	<i>P</i> -value	Observed	Expected	% difference	Z- score ¹⁴	P -value
Characters in "Les Miserables" ¹	77	2.64	2.50	5.6	3.58	0.0003	0.56	0.29	93.4	30.12	<10-4
Words in "David Copperfield"2	112	2.54	2.48	2.3	1.81	0.0703	0.31	0.29	4.8	1.67	0.0949
Dolphins ³	62	3.36	2.70	24.3	14.40	<10-4	0.53	0.37	40.8	11.59	<10^4
Political blogs ⁴	1224	2.74	2.59	5.7	23.5	<10-4	0.43	0.14	206.9	189.27	<10-4
Co-authorship ⁵	7610	7.03	5.42	29.6	64.70	<10^4	0.81	0.49	64.9	12.50	<10 ⁻⁴
Football ⁶	115	2.51	2.23	12.5	54.30	<10^4	0.60	0.28	119.2	44.68	<10-4
Power ⁷	4941	18.99	8.32	128.3	14.30	<10^4	0.93	0.73	28.5	105.10	<10^4
Airline [®]	810	3.06	2.61	17.4	3.53	0.0004	0.31	0.13	130.0	114.70	<10-4
Electronic circuits®	512	6.86	5.64	21.6	12.40	<10^4	0.81	0.63	28.6	35.96	<10^4
Protein-protein Interaction ¹⁰	1870	6.81	5.78	17.8	9.19	<10-4	0.81	0.72	13.2	18.23	<10-4
Neural ¹¹	297	2.46	2.35	4.5	3.38	0.0007	0.40	0.22	80.0	51.26	<10-4
Transcriptional regulatory ¹²	3459	3.72	3.39	9.7	3.60	0.0003	0.60	0.47	29.5	58.29	<10 ⁻⁴
Metabolic ¹³	563	8.78	6.54	34.3	18.67	<10 ⁻⁴	0.84	0.73	14.5	14.72	<10-4

¹The network of coappearances of characters in Victor Hugo's novel "Les Miserables". Nodes represent characters and edges connect any pair of characters that appear in the same chapter.

³The network of common adjective and noun adjacencies for the novel "David Copperfield" by Charles Dickens. Nodes represent the most commonly occurring adjectives and nouns in the book.

³The network of frequent associations between 62 dolphins in a community living off Doubtful Sound, New Zealand.

⁴The network of political blogs. Nodes represent blogs and edges are the links between blogs.

⁵The network of scientists posting preprints on the high-energy theory archive at www.antiv.org, 1995–1999. Nodes are authors and edges connect coauthors. ⁶The network of American football games between Division IA colleges during regular season Fall 2000. Nodes are teams and edges connect teams that contest in a game.

⁷The network of the Western States Power Grid of the United States. Nodes are power plants, stations and households, and edges are powerlines.

"The network of scheduled air line connections in United States, 2005. Nodes are airports and edges are scheduled direct flights.

⁹Electronic circuits. Nodes are electronic elements and edges are electronic connections.

¹⁰The protein-protein interaction network of the budding yeast 5. cerevisioe. Nodes are proteins and edges connect proteins that interact with each other.
¹¹The neural network for the worm C. elegans. Nodes are neurons and edges link neurons that connect.

12The transcriptional regulatory network of the budding yeast S. cerevisiae. Nodes are genes and edges connect genes that regulate one another.

13The metabolic network of the bacterium E. coll. Nodes are metabolites and edges connect metabolites that can be converted by a biochemical reaction.

¹⁶Z-score, number of standard deviations by which the observation deviates from the expectation.

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http://complexnt.blogspot.com

Introduction to Network Science

Generating Functions and Degree Distributions

The generating function (gf) for the probability distribution p_k is the polynomial

$$g(z) = \sum_{k=0}^{\infty} p_k z^k.$$

If we know gf for p_k then we can recover the values of p_k by differentiating

$$p_k = \frac{1}{k!} \frac{d^k g}{dz^k} \bigg|_{z=0}$$

Example: k = 0, 1, 2 with the respective $p_k = \frac{1}{2}, \frac{7}{16}, \frac{1}{16}$ for all k then

$$g(z) = \frac{1}{2} + \frac{7}{16}z + \frac{1}{16}z^2$$

Example: k follows Poisson distribution, i.e., $p_k = e^{-c} \frac{c^k}{k!}$

$$g(z) = e^{-c} \sum_{k=0}^{\infty} \frac{c^k}{k!} z^k = e^{c(z-1)}$$

Introduction to Network Science

Power-law distributions $p_k = Ck^{-\alpha}$, $\alpha > 0$, k > 0Reminder: *C* is calculated from normalization condition, i.e., $C = 1/\zeta(\alpha)$

Since we are interested in differentiating g(z) note that

$$\frac{\partial Li_{\alpha}(z)}{\partial z} = \frac{Li_{\alpha-1}(z)}{z}$$

Some properties of g(z)

• g(1) = 1

•
$$\langle k \rangle = g'(1), \, \langle k^2 \rangle = \left[\left(z \frac{d}{dz} \right)^2 g(z) \right]_{z=1}, \, \dots, \, \langle k^m \rangle = \left[\left(z \frac{d}{dz} \right)^m g(z) \right]_{z=1}$$

• Choose *m* integers k_i from $p_k \Rightarrow \Pr[\text{chosing particular set of values } \{k_i\}] = \prod_i p_{k_i}$

$$\pi_s = \Pr[\sum_{i=1}^m k_i = s] = \sum_{k_1=0}^\infty \cdots \sum_{k_m=0}^\infty \delta(s, \sum_i k_i) \prod_{i=0}^m p_{k_i} \Rightarrow h(z) = \sum_{s=0}^\infty \pi_s z^s = \cdots = \left(\sum_{k=0}^\infty p_k z_k\right)^m = (g(z))^m$$
drawn values add to a specific sum *s*

Introduction to Network Science

Random Graphs and Configuration Model

Degrees: 1, 1, 2, 2, 3, 3



1. Add *n* nodes

- 2. Add initial d(i) stubs to each i
- 3. Connect stubs iteratively

Problems? Total degree is even; Can create self-loops, multi-edges

Configuration Model

Multi-edges: Probability of adding an edge between i and j with degrees k_i , and k_j is

$$p_{ij} = \frac{k_i k_j}{2m - 1} \not \sim$$

in the limit we can omit -1

Probability of second edge is $(k_i - 1)(k_j - 1)/2m$

Expected number of multiedges in conf model

$$\frac{1}{2(2m)^2} \sum_{ij} k_i k_j (k_i - 1)(k_j - 1) = \frac{1}{2\langle k \rangle^2 n^2} \sum_i k_i (k_i - 1) \sum_j k_j (k_j - 1) = \frac{1}{2} \left[\frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle} \right]^2$$

Similar result for self-edges

$$\sum_{i} p_{ii} = \sum_{i} \frac{k_i(k_i - 1)}{4m} = \frac{\langle k^2 \rangle - \langle k \rangle}{2\langle k \rangle}$$

Conclusion? Expected number of multi-edges remains constant as network grows. Expected number of common neighbors

$$n_{ij} = \sum_{l} \frac{k_i k_l}{2m} \frac{k_j (k_l - 1)}{2m} = \frac{k_i k_j}{2m} \frac{\sum_{l} k_l (k_l - 1)}{n \langle k \rangle} = p_{ij} \frac{\langle k^2 \rangle - \langle k \rangle}{\langle k \rangle}$$
i is connected to *I i j* is connected to *I* given *iI*

Final Project

- Download full data set
- Choose one primary type, s.t. the number of full records is >10000 (don't choose "criminal trespass", remove "noisy" records with no time and coordinate)
- Build an undirected network with the following rules
 - Records are nodes
 - Ij is an edge iff (1) primary_type(i) = primary_type(j); (2) |time(i) time(j)|<p1; and (3) dist(location(i) location(j))<p2, where p1 and p2 are parameters
 - for location use longitude and latitude; for distance use Euclidean distance
- Degree distribution: 1) compute and plot (make sure you adjust the scales, and everything is visible); 2) is it similar to Poisson/binomial, exponential, power law or something else? 3) estimate the chance that "high impact" crime will be connected to another "high-impact" crime (use excess degree distribution and/or degree distr of neighbors)
- Importance of nodes: use various centrality indices to model importance of nodes (compute, plot, explain).
- Compute or estimate clustering coefficient of a network (explain)
- Compute modularity.
- Find clusters in the network (e.g., graclus), plot distribution of sizes
- Can you disconnect the network? What is the smallest size of separator? (e.g., metis)
- Visualize a subnetwork for 1 year (or for less than 1 year)