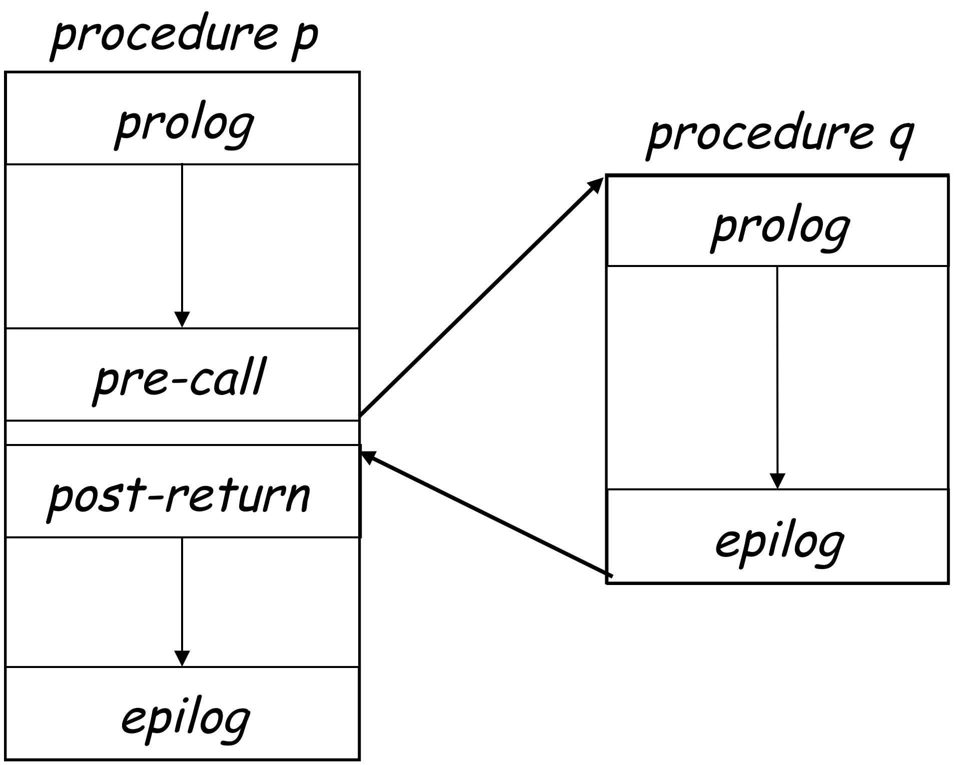
Code Shape II
Procedure Calls, Dispatch, Booleans, Relationals, & Control flow

Procedure Linkages

Standard procedure linkage



Procedure has
- standard prolog
- standard epilog

Each call involves a
- pre-call sequence
- post-return sequence

These are completely predictable from the call site ⇒ depend on the number & type of the actual parameters

Implementing Procedure Calls

If *p* calls *q*, one of them must

- Preserve register values (*caller-saves versus callee saves*)
  - → Caller-saves registers stored/restored by *p* in *p*'s AR
  - → Callee-saves registers stored/restored by *q* in *q*'s AR
- Allocate the AR
  - → Heap allocation ⇒ callee allocates its own AR
  - → Stack allocation ⇒ caller & callee cooperate to allocate AR

Space tradeoff

- Pre-call & post-return occur on every call
- Prolog & epilog occur once per procedure
- More calls than procedures
  - → Moving operations into prolog/epilog saves space

Implementing Procedure Calls

If *p* calls *q*, one of them must

- Preserve register values (caller-saves versus callee saves)

If space is an issue

- Moving code to prolog & epilog saves space
- As register sets grow, save/restore code does, too
  - → Each saved register costs 2 operations
  - → Can use a library routine to save/restore
    - ♦ Pass it a mask to determine actions & pointer to space
    - ♦ Hardware support for save/restore or storeM/loadM

Can decouple who saves from what is saved

Implementing Procedure Calls

Evaluating parameters

- Call by reference $\Rightarrow$ evaluate parameter to an lvalue
- Call by value $\Rightarrow$ evaluate parameter to an rvalue & store it

Aggregates (structs), arrays, & strings are usually c-b-r

- Language definition issues
- Alternatives
  - → Small structures can be passed in registers
  - → Can pass large c-b-v objects c-b-r and copy on modification

Procedure-valued parameters

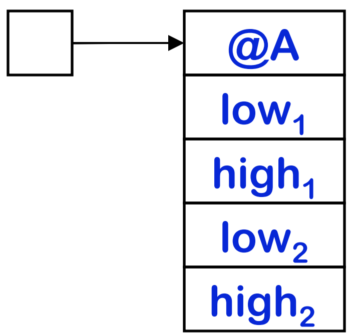
- Must pass starting address of procedure

Implementing Procedure Calls

What about arrays as actual parameters?

Whole arrays, as call-by-reference parameters

- Callee needs dimension information
  - → Builds a descriptor called a *dope vector*
- Store the values in the calling sequence
- Pass the address of the dope vector in the parameter slot
- Generate complete address polynomial at each reference



Implementing Procedure Calls

What about A[12] as an actual parameter?

If corresponding parameter is a scalar, it's easy

- Pass the address or value, as needed

What if corresponding parameter is an array?

- See previous slide

Implementing Procedure Calls

What about a string-valued argument?

- Call by reference ⇒ pass a pointer to the start of the string
  - → Works with either length/contents or null-terminated string
- Call by value ⇒ copy the string & pass it
  - → Can store it in caller's AR or callee's AR
  - → Can pass by reference & have callee copy it if necessary …

Pointer of string serves as "descriptor" for the string, stored in the appropriate location (register or slot in the AR)

Implementing Procedure Calls

What about a structure-valued parameter?

- Again, pass a handle
- Call by reference $\Rightarrow$ descriptor (pointer) refers to original
- Call by value $\Rightarrow$ create copy & pass its descriptor
  - → Can allocate it in either caller's AR or callee's AR
  - → Can pass by reference & have callee copy it if necessary …

If it is actually an array of structures, then use a dope vector

What About Calls in an OOL (Dispatch)?

In an OOL, most calls are indirect calls

- Compiled code does not contain address of callee
  - → Finds it by indirection through class' method table
  - → Required to make subclass calls find right methods
  - → Code compiled in class *C* cannot know of subclass methods that override methods in *C* and *C*'s superclasses
- In the general case, need dynamic dispatch
  - → Map method name to a search key
  - → Perform a run-time search through hierarchy
    - ♦ Start with object's class, search for 1st occurrence of key
    - ♦ This can be expensive
  - → Use a method cache to speed search
    - ♦ Cache holds < ***key, class, method pointer*** >

How big?
Bigger ⇒ more hits & longer search
Smaller ⇒ fewer hits, faster search

What About Calls in an OOL (Dispatch)?

Improvements are possible in special cases

- If class has no subclasses, can generate direct call
  - → Class structure must be static or class must be **FINAL**
- If class structure is static
  - → Can generate complete method table for each class
  - → Single indirection through class pointer *(1 or 2 operations)*
  - → Keeps overhead at a low level
- If class structure changes infrequently
  - → Build complete method tables at run time
  - → Initialization & any time class structure changes

What About Calls in an OOL (Dispatch)?

Unusual issues in OOL call

- Need to pass receiver's object record as (1st) parameter
  - → Becomes self or this
- Method needs access to its class
  - → Object record has static pointer to superclass, and so on …
- Method is a full-fledged procedure
  - → It still needs an AR …
  - → Can often stack allocate them *(**HotSpot does** …)*

Boolean & Relational Values

How should the compiler represent them?

- Answer depends on the target machine

Two classic approaches

- Numerical representation
- Positional (implicit) representation

Correct choice depends on both context and ISA

Boolean & Relational Values

Numerical representation

- Assign values to TRUE and FALSE
- Use hardware AND, OR, and NOT operations
- Use comparison to get a boolean from a relational expression

Examples

x < y	*becomes*	cmp_LT $r_x, r_y \Rightarrow r_1$
if (x < y) then $stmt_1$ else $stmt_2$	*becomes*	cmp_LT $r_x, r_y \Rightarrow r_1$ cbr r1 $\rightarrow$ _$stmt_1$,_$stmt_2$

Boolean & Relational Values

What if the ISA uses a condition code?

- Must use a conditional branch to interpret result of compare
- Necessitates branches in the evaluation

Example:

x < y ***becomes***

	cmp	$r_x, r_y \Rightarrow cc_1$
	cbr_LT	$cc_1 \rightarrow L_T, L_F$
L_T:	**loadI**	$1 \Rightarrow r_2$
	br	$\rightarrow L_E$
L_F:	**loadI**	$0 \Rightarrow r_2$
L_E:	**…other stmts…**	

This "positional representation" is much more complex

Boolean & Relational Values

What if the ISA uses a condition code?

- Must use a conditional branch to interpret result of compare
- Necessitates branches in the evaluation

Example:

x < y *becomes*

```
        cmp      rx, ry ⇒ cc1
        cbr_LT   cc1 → LT, LF
LT:     loadI    1 ⇒ r2
        br       → LE
LF:     loadI    0 ⇒ r2
LE:     ...other stmts...
```

Condition codes

- are an architect's hack
- allow ISA to avoid some comparisons
- complicates code for simple cases

This "positional representation" is much more complex

Boolean & Relational Values

The last example actually encodes result in the PC

If result is used to control an operation, this may be enough

Example
if (x < y)
then a ← c + d
else a ← e + f

Variations on the ILOC Branch Structure					
Straight Condition Codes			*Boolean Compares*		
	comp	$r_x, r_y \Rightarrow cc_1$		cmp_LT	$r_x, r_y \Rightarrow r_1$
	cbr_LT	$cc_1 \rightarrow L_1, L_2$		cbr	$r_1 \rightarrow L_1, L_2$
L_1:	add	$r_c, r_d \Rightarrow r_a$	L_1:	add	$r_c, r_d \Rightarrow r_a$
	br	$\rightarrow L_{OUT}$		br	$\rightarrow L_{OUT}$
L_2:	add	$r_e, r_f \Rightarrow r_a$	L_2:	add	$r_e, r_f \Rightarrow r_a$
	br	$\rightarrow L_{OUT}$		br	$\rightarrow L_{OUT}$
L_{OUT}:	nop		L_{OUT}:	nop	

Condition code version does not directly produce (x < y)

Boolean version does

Still, there is no significant difference in the code produced

Boolean & Relational Values

Conditional move & predication both simplify this code

Example
if (x < y)
then a ← c + d
else a ← e + f

Other Architectural Variations				
Conditional Move		***Predicated Execution***		
comp	$r_x, r_y \Rightarrow cc_1$		cmp_LT	$r_x, r_y \Rightarrow r_1$
add	$r_c, r_d \Rightarrow r_1$	(r_1)?	add	$r_c, r_d \Rightarrow r_a$
add	$r_e, r_f \Rightarrow r_2$	$(\neg r_1)$?	add	$r_e, r_f \Rightarrow r_a$
i2i_<	$cc_1, r_1, r_2 \Rightarrow r_a$			

Both versions avoid the branches

Both are shorter than CCs or Boolean-valued compare

Are they better?

Boolean & Relational Values

Consider the assignment $x \leftarrow a < b \land c < d$

VARIATIONS ON THE ILOC BRANCH STRUCTURE	
Straight Condition Codes	***Boolean Compare***
comp $r_a, r_b \Rightarrow cc_1$ cbr_LT $cc_1 \rightarrow L_1, L_2$ L_1: comp $r_c, r_d \Rightarrow cc_2$ cbr_LT $cc_2 \rightarrow L_3, L_2$ L_2: loadI $0 \Rightarrow r_x$ br $\rightarrow L_{OUT}$ L_3: loadI $1 \Rightarrow r_x$ br $\rightarrow L_{OUT}$ L_{OUT}: nop	cmp_LT $r_a, r_b \Rightarrow r_1$ cmp_LT $r_c, r_d \Rightarrow r_2$ and $r_1, r_2 \Rightarrow r_x$

Here, the boolean compare produces much better code

Boolean & Relational Values

Conditional move & predication help here, too

$x \leftarrow a < b \wedge c < d$

OTHER ARCHITECTURAL VARIATIONS					
Conditional Move			***Predicated Execution***		
comp	r_a, r_b	$\Rightarrow cc_1$	cmp_LT	$r_a, r_b \Rightarrow r_1$	
i2i_<	cc_1, r_T, r_F	$\Rightarrow r_1$	cmp_LT	$r_c, r_d \Rightarrow r_2$	
comp	r_c, r_d	$\Rightarrow cc_2$	and	$r_1, r_2 \Rightarrow r_x$	
i2i_<	cc_2, r_T, r_F	$\Rightarrow r_2$			
and	r_1, r_2	$\Rightarrow r_x$			

Conditional move is worse than Boolean compares

Predication is identical to Boolean compares

Context & hardware determine the appropriate choice

Control Flow

If-then-else

- Follow model for evaluating relationals & booleans with branches

Branching versus predication *(e.g., IA-64)*

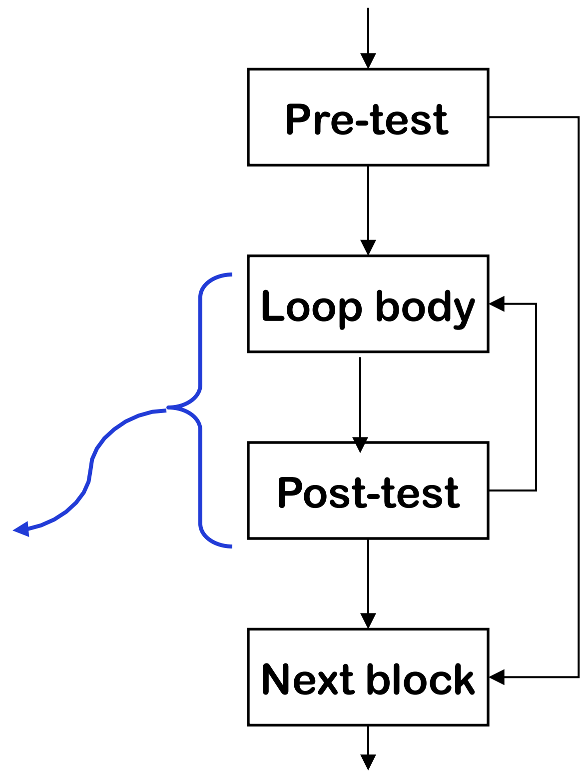
- Frequency of execution
  - → Uneven distribution ⇒ do what it takes to speed common case
- Amount of code in each case
  - → Unequal amounts means predication may waste issue slots
- Control flow inside the construct
  - → Any branching activity within the case base complicates the predicates and makes branches attractive

Control Flow

Loops

- Evaluate condition before loop (if needed)
- Evaluate condition after loop
- Branch back to the top (if needed)

Merges test with last block of loop body



while, for, do, & until all fit this basic model

Loop Implementation Code

for (i = 1; i< 100; i++) { *body* }

next statement

	loadI	$1 \Rightarrow r_1$	Initialization
	loadI	$1 \Rightarrow r_2$	
	loadI	$100 \Rightarrow r_3$	Pre-test
	cmp_GE	$r_1, r_3 \Rightarrow r_4$	
	cbr	$r_4 \Rightarrow L_2, L_1$	
L_1:	*body*		
	add	$r_1, r_2 \Rightarrow r_1$	
	cmp_LT	$r_1, r_3 \Rightarrow r_5$	Post-test
	cbr	$r_5 \Rightarrow L_1, L_2$	
L_2:	*next statement*		

Break statements

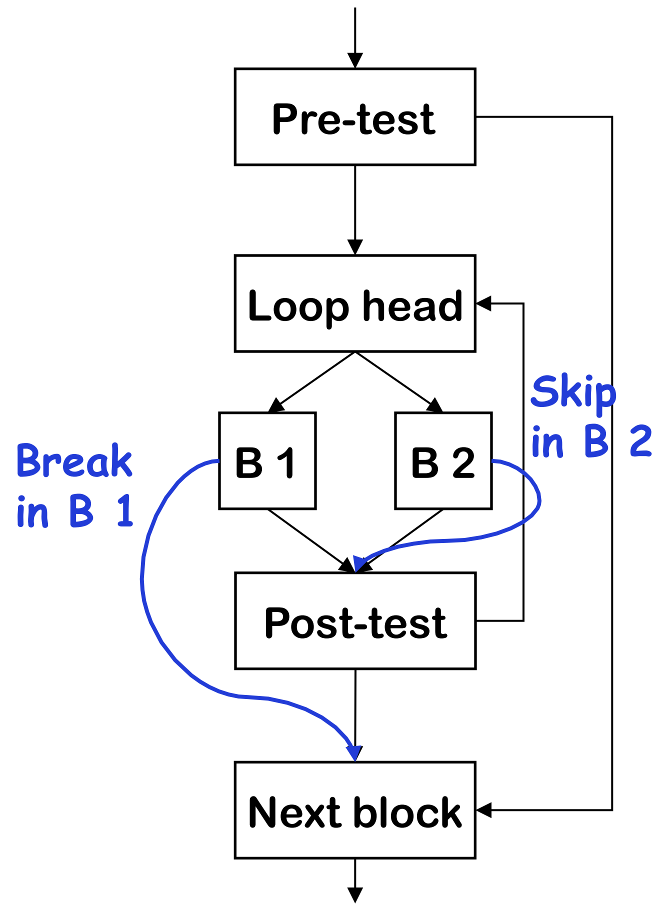
Many modern programming languages include a break

- Exits from the innermost control-flow statement
  - → Out of the innermost loop
  - → Out of a case statement

Translates into a jump

- Targets statement outside control-flow construct
- Creates multiple-exit construct
- Skip in loop goes to next iteration

Only make sense if loop has > 1 block



Control Flow

Case Statements

1 Evaluate the controlling expression

2 Branch to the selected case

3 Execute the code for that case

4 Branch to the statement after the case

Parts 1, 3, & 4 are well understood, part 2 is the key

Control Flow

Case Statements

1 Evaluate the controlling expression

2 Branch to the selected case

3 Execute the code for that case

4 Branch to the statement after the case (*use break*)

Parts 1, 3, & 4 are well understood, part 2 is the key

Surprisingly many compilers do this for all cases!

Strategies

- Linear search (nested if-then-else constructs)
- Build a table of case expressions & binary search it
- Directly compute an address (requires dense case set)