

Context-sensitive Analysis

Beyond Syntax



There is a level of correctness that is deeper than grammar

fie(a,b,c,d) int a, b, c, d; { ... } fee() { int f[3],g[0],h,i,j,k; char *p; fie(h,i,"ab",j, k); k = f * i + j;h = g[17];printf("< $%s, %s > . \n", p, q$); p = 10:}

What is wrong with this program?

Beyond Syntax



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What is wrong with this program?

- declared g[0], used g[17]
- wrong number of args to fie()
- "ab" is not an <u>int</u>
- wrong dimension on use of f
- undeclared variable q
- 10 is not a character string

All of these are

"deeper than syntax"

To generate code, we need to understand its meaning !



- To generate code, the compiler needs to answer many questions
- Is "x" a scalar, an array, or a function? Is "x" declared?
- Are there names that are not declared? Declared but not used?
- Which declaration of "x" does each use reference?
- Is "x" defined before it is used?
- Is the expression "x * y + z" type-consistent?

These are beyond a Context Free Grammar



- To generate code, the compiler needs to answer many questions
- In "*a*[i,j,k]", does *a* have three dimensions?
- Where can "z" be stored? (*register, local, global, heap, static*)
- How many arguments does "fie()" take?
- Does "*p" reference the result of a "malloc()" ?
- Do "p" & "q" refer to the same memory location?

These are beyond a Context Free Grammar

Beyond Syntax



These questions are part of context-sensitive analysis

- Questions & answers involve non-local information
- Answers may involve computation

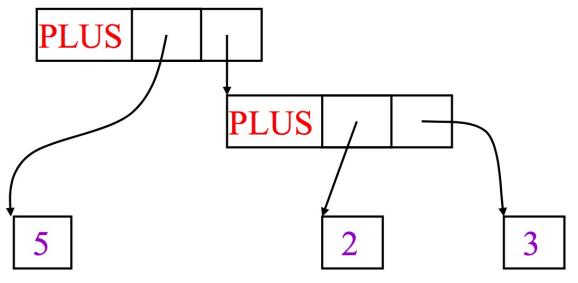
How can we answer these questions?

- Use formal methods
 - \rightarrow Attribute grammars?
 - Also known as attributed CFG or syntax-directed definitions
- Use *ad-hoc* techniques
 - \rightarrow Symbol tables
 - \rightarrow *Ad-hoc* code

(action routines)

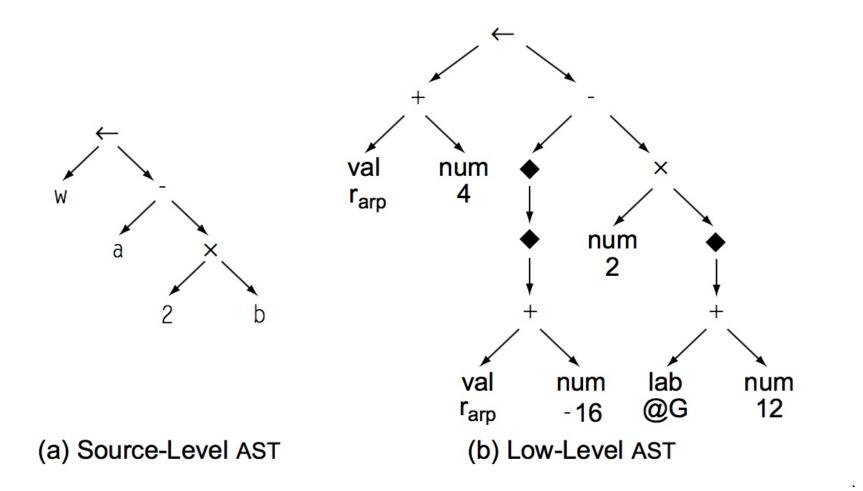
In scanning & parsing, formalism won; different story here.





- Also captures the nesting structure
- But <u>abstracts</u> from the concrete syntax
 → more compact and easier to use
- An important data structure in a compiler





ASTs and Parsing



- AST can be built doing a bottom-up parse
- The construction procedure is rather simple
 - Create appropriate tree nodes for each element in the grammar
 - for each node, carry sufficient essential information of program fragments it represents

What type of trees they have built by using CUP?

Read the code of the skeleton to understand what they are operating on

Inheritance Graph



