



Context-sensitive Analysis

Part IV

Ad-hoc syntax-directed translation,
Symbol Tables, and Types

Quiz



Name two differences between attribute grammars and ad-hoc syntax directed translation techniques?



Example: Processing C Declarations

<i>DeclarationList</i>	→	<i>DeclarationList Declaration</i>
		<i>Declaration</i>
<i>Declaration</i>	→	<i>SpecifierList InitDeclaratorList ;</i>
<i>SpecifierList</i>	→	<i>Specifier SpecifierList</i>
		<i>Specifier</i>
<i>Specifier</i>	→	<i>StorageClass</i>
		<i>TypeSpecifier</i>
<i>StorageClass</i>	→	auto
		static
		extern
		register
<i>TypeSpecifier</i>	→	void
		char
		short
		int
		long
		signed
		unsigned
		float
		double



Example: Processing C Declarations

DeclarationList → *DeclarationList Declaration*
|
Declaration

Declaration → *SpecifierList InitDeclaratorList ;*

SpecifierList → *Specifier SpecifierList*
|
Specifier

Specifier → *StorageClass*
|
TypeSpecifier

StorageClass → auto
|
static
|
extern
|
register

TypeSpecifier → void
|
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Example: Processing C Declarations

<i>InitDeclaratorList</i>	→	<i>InitDeclaratorList</i> , <i>InitDeclarator</i>
		<i>InitDeclarator</i>

<i>InitDeclarator</i>	→	<i>Declarator</i> = <i>Initializer</i>
		<i>Declarator</i>

<i>Declarator</i>	→	<i>Pointer</i> <i>DirectDeclarator</i>
		<i>DirectDeclarator</i>

<i>Pointer</i>	→	*
		* <i>Pointer</i>

<i>DirectDeclarator</i>	→	<i>ident</i>
		(<i>Declarator</i>)
		<i>DirectDeclarator</i> ()
		<i>DirectDeclarator</i> (<i>ParameterTypeList</i>)
		<i>DirectDeclarator</i> (<i>IdentifierList</i>)
		<i>DirectDeclarator</i> []
		<i>DirectDeclarator</i> [<i>ConstantExpr</i>]

Example: Processing C Declarations

InitDeclaratorList → *InitDeclaratorList* , *InitDeclarator*
| *InitDeclarator*

InitDeclarator → *Declarator* = *Initializer*
| *Declarator*

Declarator → *Pointer* *DirectDeclarator*
| *DirectDeclarator*

Pointer → *
| * *Pointer*

DirectDeclarator → *ident*
| (*Declarator*)
| *DirectDeclarator* ()
| *DirectDeclarator* (*ParameterTypeList*)
| *DirectDeclarator* (*IdentifierList*)
| *DirectDeclarator* []
| *DirectDeclarator* [*ConstantExpr*]

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Building a Symbol Table

- Enter declaration information as processed
- At end of declaration syntax, do some post processing
- Use table to check errors as parsing progresses

assumes table
is global



Symbol Table: Typical Uses

- Define before use → lookup on reference
- Dimension, type, ... → check as encountered
- Type checking of expression → bottom-up walk



Symbol Table: Typical Uses

- Procedure interfaces are harder
 - Build a representation for parameter list & types
 - Create list of sites to check



Is This Really "Ad-hoc" ?

Relationship between practice and attribute grammars

Similarities

- Both rules & actions associated with productions
- Application order determined by tools, not author



Is This Really "Ad-hoc" ?

Relationship between practice and attribute grammars

Differences

- Actions applied as a unit; not true for AG rules
- Anything goes in *ad-hoc* actions; AG rules are functional



Making Ad-hoc SDT Work

How do we fit this into an LR(1) parser?

```
stack.push(INVALID);
stack.push( $s_0$ ); // initial state
token = scanner.next_token();
loop forever {
    s = stack.top();
    if ( ACTION[s,token] == "reduce  $A \rightarrow \beta$ " ) then {
        stack.popnum( $2 * |\beta|$ ); // pop  $2 * |\beta|$  symbols
        s = stack.top();
        stack.push(A); // push A
        stack.push(GOTO[s,A]); // push next state
    }
    else if ( ACTION[s,token] == "shift  $s_i$ " ) then {
        stack.push(token); stack.push( $s_i$ );
        token ← scanner.next_token();
    }
    else if ( ACTION[s,token] == "accept"
              & token == EOF )
        then break;
    else throw a syntax error;
}
report success;
```

From an earlier lecture



```
stack.push(INVALID);
stack.push(NULL);
stack.push(s0); // initial state
token = scanner.next_token();
loop forever {
    s = stack.top();
    if ( ACTION[s,token] == "reduce A→β" ) then {

        /* insert case statement here */

        stack.popnum(3*|β|); // pop 3*|β| symbols
        s = stack.top();
        stack.push($$); // push result
        stack.push(A); // push A
        stack.push(GOTO[s,A]); // push next state
    }
    else if ( ACTION[s,token] == "shift si" ) then {
        stack.push(attr); stack.push(token);
        stack.push(si);
        token ← scanner.next_token();
    }
    else if ( ACTION[s,token] == "accept"
              & token == EOF )
        then break;
    else throw a syntax error;
}
report success;
```

To add yacc-like actions

- Stack has 3 items per symbol rather than 2 (3rd is \$\$)



```
stack.push(INVALID);
stack.push(NULL);
stack.push( $s_0$ ); // initial state
token = scanner.next_token();
loop forever {
    s = stack.top();
    if ( ACTION[s,token] == "reduce  $A \rightarrow \beta$ " ) then {
        /* insert case statement here */
        stack.popnum( $3 * |\beta|$ ); // pop  $3 * |\beta|$  symbols
        s = stack.top();
        stack.push($$); // push result
        stack.push(A); // push A
        stack.push(GOTO[s,A]); // push next state
    }
    else if ( ACTION[s,token] == "shift  $s_i$ " ) then {
        stack.push(attr); stack.push(token);
        stack.push( $s_i$ );
        token ← scanner.next_token();
    }
    else if ( ACTION[s,token] == "accept"
              & token == EOF )
        then break;
    else throw a syntax error;
}
report success;
```

- Add case statement to the reduction processing section
→ Case switches on production number

```
stack.push(INVALID);
stack.push(NULL);
stack.push( $s_0$ ); // initial state
token = scanner.next_token();
loop forever {
    s = stack.top();
    if ( ACTION[s,token] == "reduce  $A \rightarrow \beta$ " ) then {
        /* insert case statement here */ ←
        stack.popnum( $3 * |\beta|$ ); // pop  $3 * |\beta|$  symbols
        s = stack.top();
        stack.push($$); // push result
        stack.push(A); // push A
        stack.push(GOTO[s,A]); // push next state
    }
    else if ( ACTION[s,token] == "shift  $s_i$ " ) then {
        stack.push(attr); stack.push(token);
        stack.push( $s_i$ );
        token ← scanner.next_token();
    }
    else if ( ACTION[s,token] == "accept"
              & token == EOF )
        then break;
    else throw a syntax error;
}
report success;
```

→ Each case clause holds the code snippet for that production

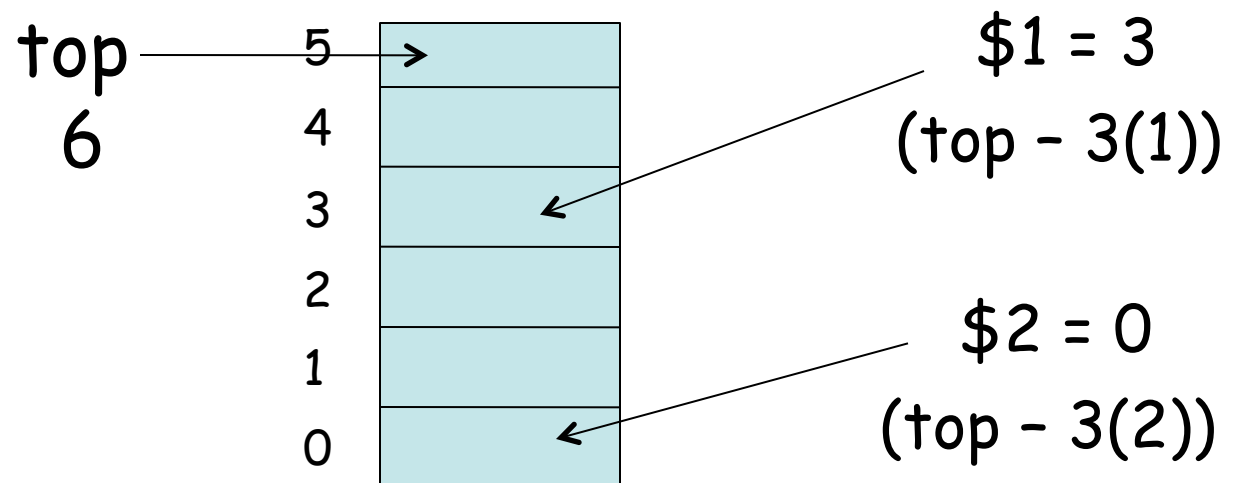
→ Substitute appropriate names for \$\$, \$1, \$2, ...



Making Ad-hoc SDT Work

How do we fit this into an LR(1) parser?

- Need a naming scheme to access them
 - $\$n$ translates into stack location ($\text{top} - 3n$)
- Need to sequence rule applications
 - On every reduce action, perform the action rule
 - Add a giant case statement to the parser





Making Ad-hoc SDT Work

What about a rule that must work in mid-production?

- Can transform the grammar
 - Split it into two parts at the point where rule must go
 - Apply the rule on reduction to the appropriate part
- Can also handle reductions on shift actions
 - Add a production to create a reduction
 - Was: $fee \rightarrow \underline{fum}$
 - Make it: $fee \rightarrow fie \rightarrow \underline{fum}$
and tie the action to this new reduction

Together, these let us apply rule at any point in the parse



Alternative Strategy

What if you need to perform actions that do not fit well into the Ad-hoc Syntax-Directed Translation framework?

- Build the abstract syntax tree using SDT
- Perform the actions during one or more treewalks
 - In an OOL, think of this problem as a classic application of the visitor pattern
 - Perform arbitrary computation in treewalk order
 - Make multiple passes if necessary

Again, a competent junior or senior CS major would derive this solution after a couple of minutes of thought.