CISC 320 Midterm Exam

Friday, March 26, 2010

Name:

There are 23 questions. The first 18 questions count 3 points each.

True or False: write T or F in the blank.

Multiple Choice: write the letter of the best choice in the blank.

Very Short Answer: write the very short answer in the blank.

Then there is a 6 point question and four 10 point questions, for a total of 100 points.

- 1. True or false: There cannot be a comparison based algorithm for finding the maximum and $\overline{\text{minimum of an array of n numbers using fewer than 5n/4 comparisons.}$
- 2. For the problem of finding the maximum and the second largest of an array of n numbers, the best algorithm we discussed uses about comparisons.
 - (a) n 1
 - (b) $n \lg(n)$
 - (c) $n + \lg(n)$
 - (d) 3n/2
- 3. Algorithm A manipulates an array of n items. It requires $\Theta(n^2)$ time. For arrays of 1000 items algorithm A takes about 1 millisecond. On an array of 1,000,000 items how long do you expect algorithm A to take?
 - (a) about 17 milliseconds
 - (b) about 17 seconds
 - (c) about 17 minutes.
 - (d) about 17 hours.
- 4. True or False: Nobody knows whether or not there might be a comparison based sorting method that uses $\Omega(n)$ comparisons to sort an array of *n* numbers.
- 5. True or False: Nobody knows whether or not there might be a comparison based sorting method that uses O(n) comparisons to sort an array of n numbers.
- 6. True or False: The recurrence T(n) = T(n/2) + O(1), for $n \ge 2$ with T(1) = 1 has the solution $\overline{T(n)} = O(\lg(n))$.

7. Professor Moriarty has four algorithms whose runtime functions (call them T_2, T_3, T_4, T_5) satisfy the four recurrences

$$T_k(n) = \begin{cases} 1, & \text{for } n = 1, \\ kT_k(n/k) + O(n), & \text{for } n > 1, \end{cases}, \text{ for } k = 2, 3, 4, 5.$$

Which best describes the situation?

- (a) $T_k(n) = O(n)$, for each k in 2..5.
- (b) $T_k(n) = O(n \lg(n))$, for each k in 2..5.
- (c) $T_k(n) = O(n^{\lg(5)})$, for some k in 2..5.
- (d) none of the above
- 8. Devious Professor Moriarty has developed four additional algorithms whose runtime functions $\overline{(we')}$ also call them T_2, T_3, T_4, T_5 , satisfy the four recurrences,

$$T_k(n) = \begin{cases} 1, & \text{for } n = 1, \\ kT_k(n/2) + O(n), & \text{for } n > 1, \end{cases}, \text{ for } k = 2, 3, 4, 5.$$

Which best describes the situation?

- (a) $T_k(n) = O(n)$, for each k in 2..5.
- (b) $T_k(n) = O(n \lg(n))$, for each k in 2..5.
- (c) $T_k(n) = O(n^{\lg(5)})$, for some k in 2..5.
- (d) none of the above
- 9. Which is true of the recurrence

$$T(n) = \begin{cases} 1, & \text{for } n = 1, \\ 7T(n/2) + O(n^2), & \text{for } n > 1 \end{cases}$$
?

- (a) T(n) describes the runtime of a matrix multiplication algorithm.
- (b) $T(n) = O(n^{\lg(7)}).$
- (c) $T(n) = O(n^{2.81}).$
- (d) all of the above

10. Which recurrence relation describes the runtime of FFT (the fast Fourier transform).

- (a) T(n) = 2T(n/2) + O(n)
- (b) T(n) = 2T(n/2) + n/2T(2)
- (c) both of the above
- (d) none of the above
- 11. When the FFT is called on $n = 2^k$ points, a primitive root of unity, ω , is used. This ω must be a primitive *m*th root of unity, for which *m*?
 - (a) m = k
 - (b) m = n
 - (c) m = n/2

- (d) none of the above
- 12. When the FFT is called on *n* points a primitive root of unity, ω , is used. For certain X and Y, the two recursive calls work with X of the points and with the Y-th power of ω .
 - (a) X = n/2 points, Y = 2 is exponent of ω
 - (b) X = n points, Y = 1/2 is exponent of ω
 - (c) X = n/2 points, Y = 1/2 is exponent of ω
 - (d) X = n points, Y = 2 is exponent of ω
- 13. True or false: The recurrence $T(n) = 2T(n/2) + O(n^2)$ describes the worst case number of comparisons used in mergesort.
- 14. True or false: The leaves of a full binary tree are all on the same level.
- 15. A graph is a DAG if, with respect to it's depth first search forest, it has no edges.
 - (a) tree
 - (b) forward
 - (c) back
 - (d) cross
- 16. RSA public key encryption is hard to crack, because
 - (a) Modular arithmetic is a total mystery
 - (b) testing for primality is difficult
 - (c) factoring large composite numbers is difficult
 - (d) The Fermat test is fooled by some non-primes called Carmichael numbers.
- 17. Let n be the number of days since the beginning of the earth's rotation until your date of birth. What is $3^n \mod 2$?
- 18. What is $2^{2^{2^2}} \mod 7$?

19. (6 points)

(a) A directed acyclic graph can be linearized. Explain what this means.

(b) Draw a linearization of this graph: $V = \{A, B, C, D\}$, nbr[A] = C, D; nbr[B] = A, D; nbr[C] = D; nbr[D] = null.

(c) How may a linearization of a DAG be constructed from the results of depth first search that computes previsit and postvisit times for each node?

20. (10 points)

(a) Using modmul(a, b, m) which computes $a \times b \mod m$, write modexp(a, e, m), which computes $a^e \mod m$, for n bit numbers a, e, m. Your program should use no more than 2n calls to modmul, that is, no more than 2n multiplications or squarings mod m.

- (b) Up to big-O, What is the runtime of your modexp?
- 21. (5 points)
 - (a) Write mergesort(a[0..n-1]). You may use (without defining it) merge(a[0..i-1], b[0..j-1]), which merges the two sorted array segments a and b, putting the result in c[0..i + j 1].

22. (10 points)

(a) Let $f(x) = \sum_{i=0}^{n-1} f_i x^i$ be a *n* term polynomial. Explain how f(a) and f(-a) can be computed by combining $f_e(a^2)$ and $f_o(a^2)$, where f_e and f_o are the even part and the odd part of f, respectively.

- (b) When ω is an *n*-th root of unity, what power of ω is equal to $-\omega$?
- 23. (10 points)
 - (a) Write explore(G, v), which does a depth first search of the graph G = (V, E) beginning at vertex v. You may use the adjacency list representation, with nbr[u] being the list of neighbors of u in G.

- (b) Up to big-O, what is the run time of explore?
- 24. Have a good spring break.